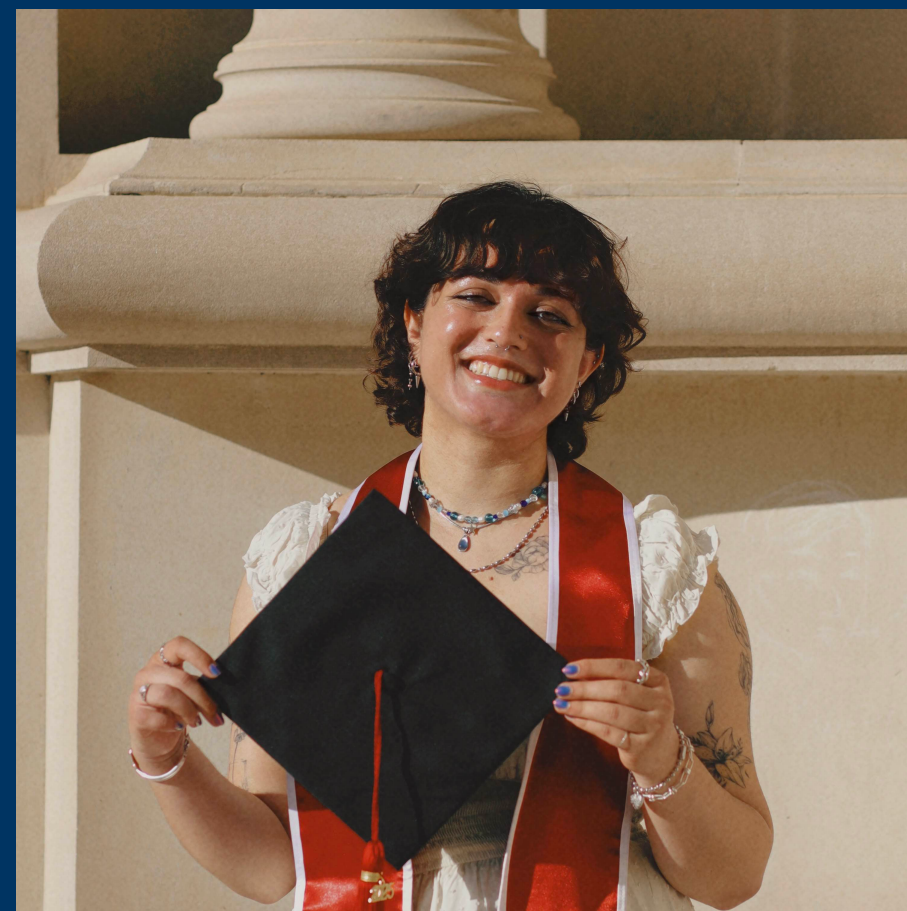


Algebra is Half the Battle: Verifying Presentations for Graded Unipotent Chevalley Groups

ITP 2025



Eric Wang



Arohee Bhoja



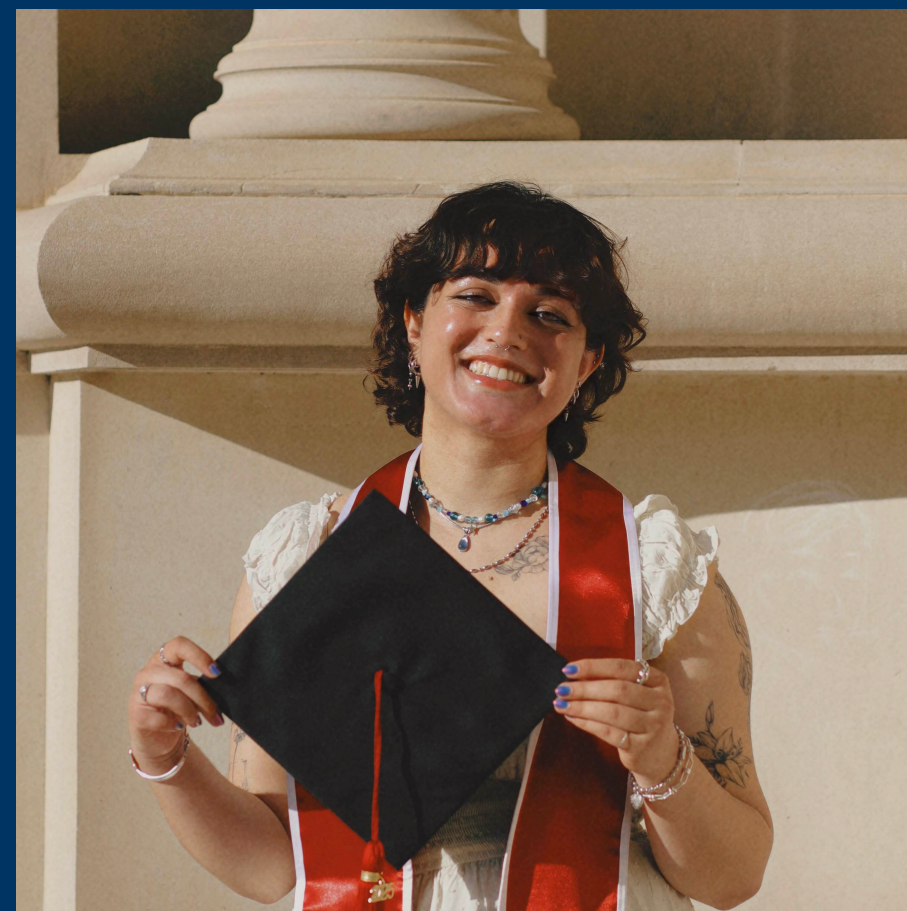
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Noah Singer



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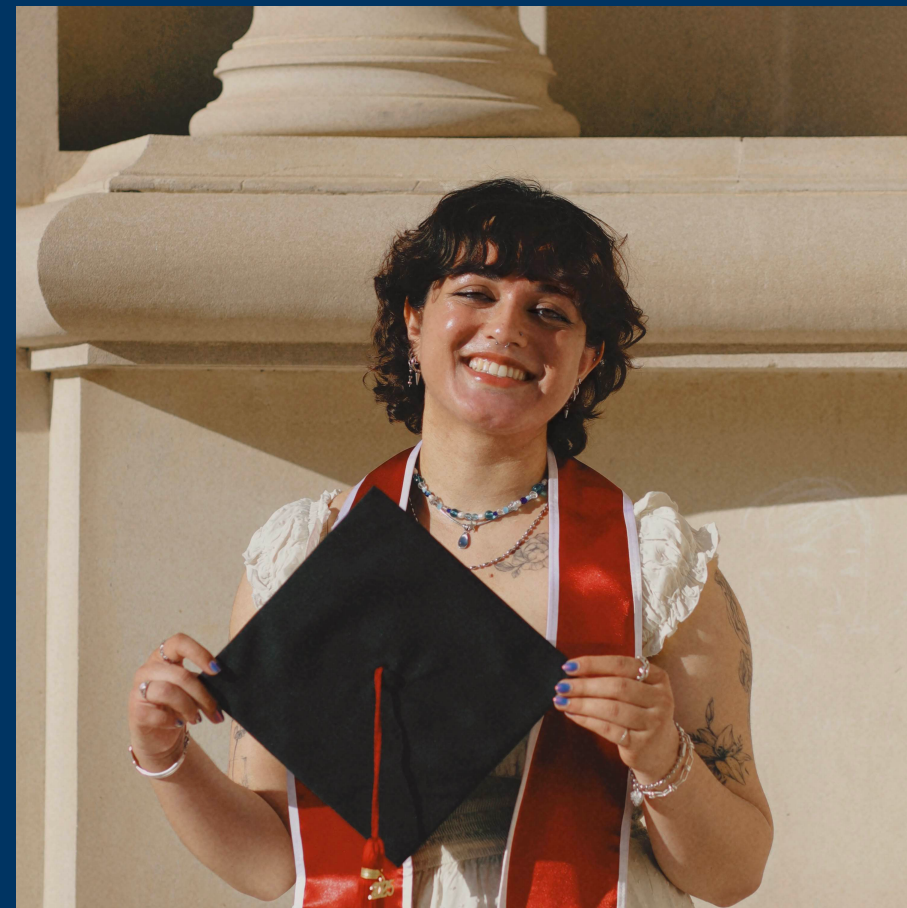


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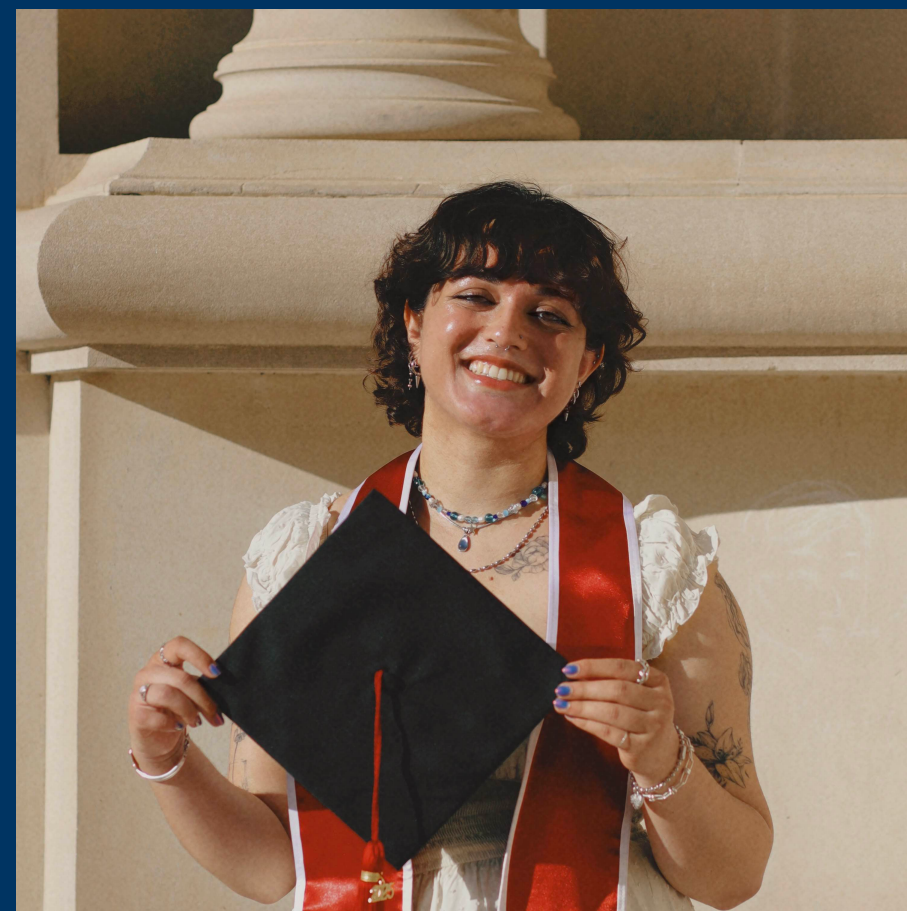
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Undergraduates
(Wrote the proofs)



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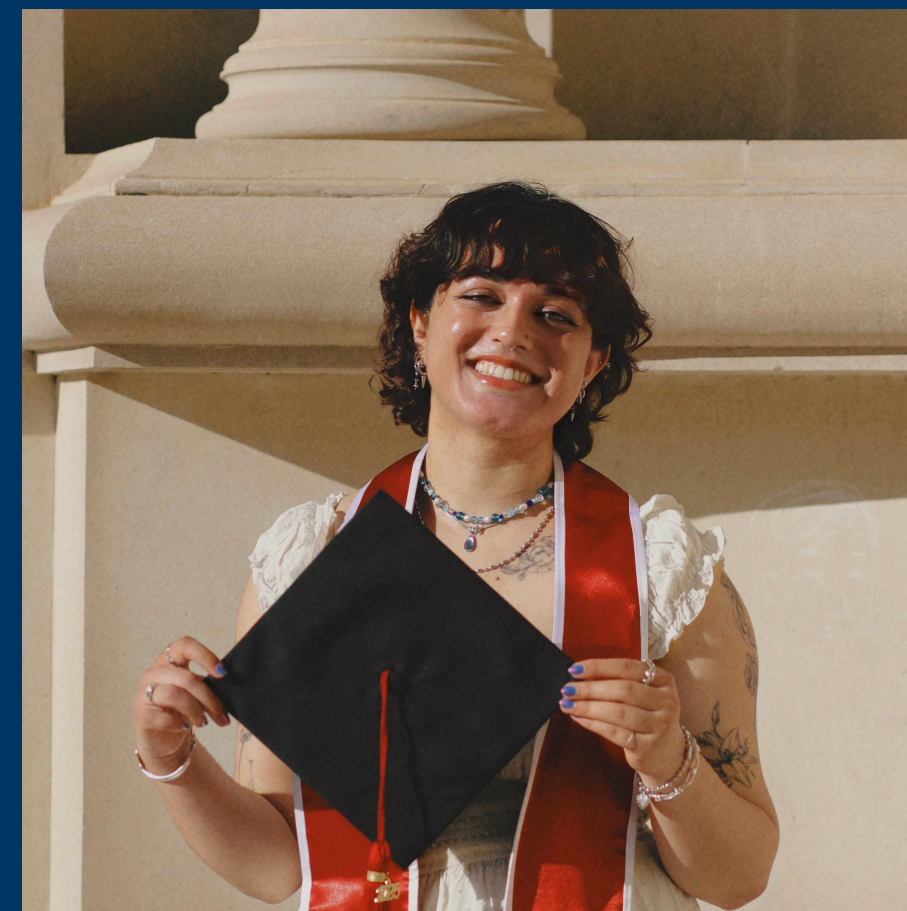


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PhD student
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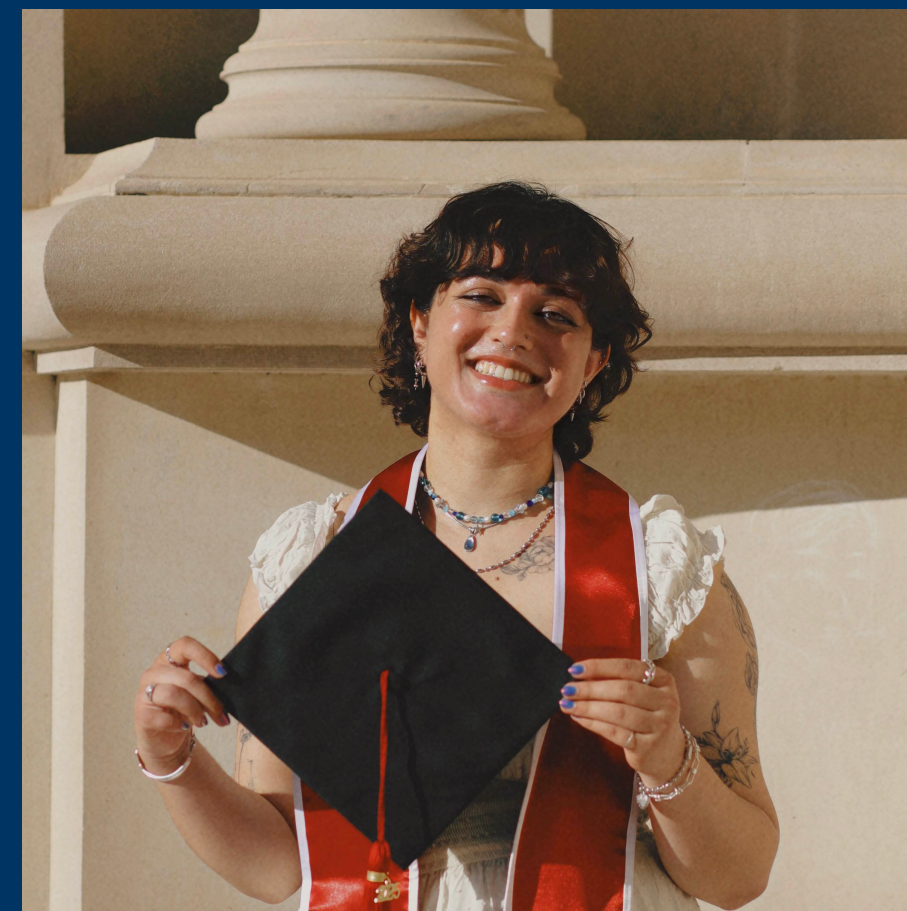
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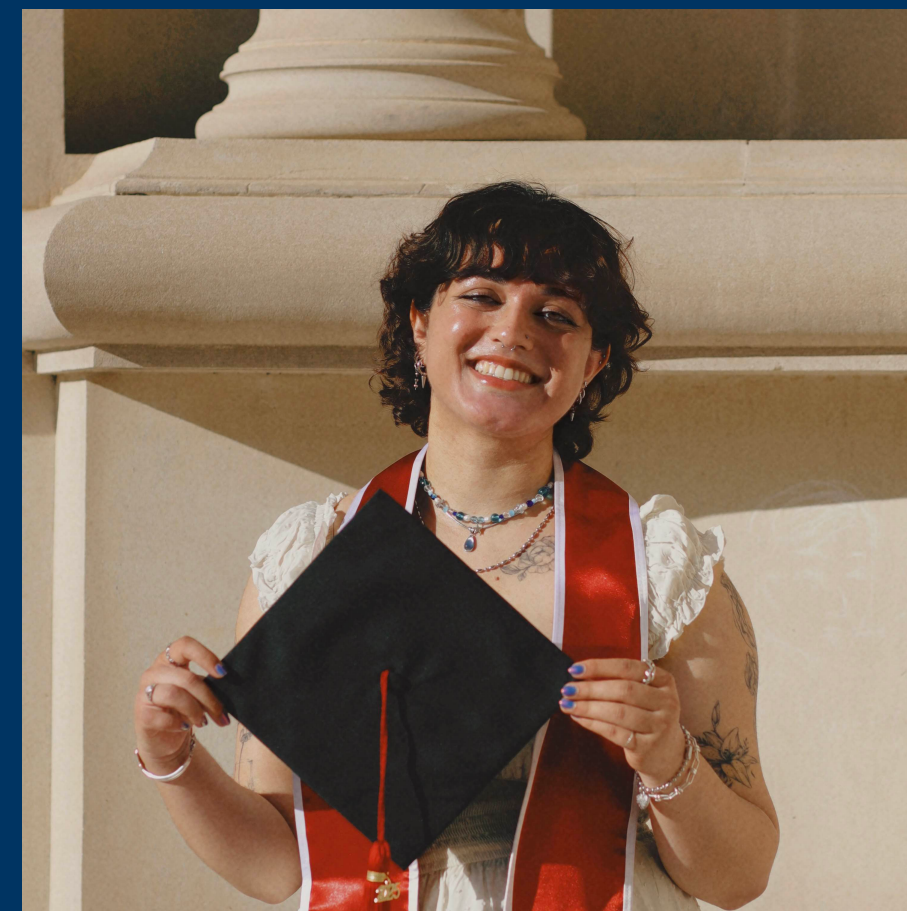
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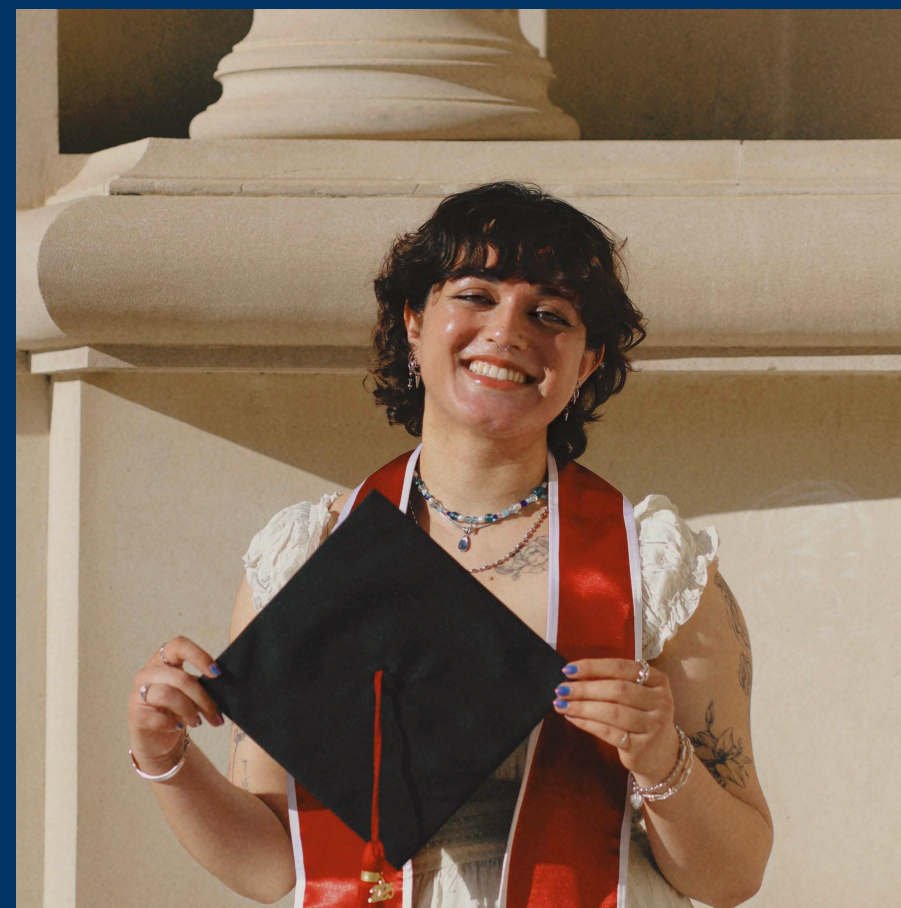
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Algebra is Half the Battle: Verifying Presentations for Graded Unipotent Chevalley Groups

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- Some type isomorphisms are better than others

$$\alpha \mid \beta \mid \alpha + \beta$$

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- Automation in need of automation `lin_id_inv_thms`

Warning!



I know about the math we formalized, but not much about the surrounding research context

What we proved

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We proved in Lean that three specific groups can be defined using a canonically smaller set of equations than previously known.

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More specifically, we showed that the A_3 , B_3 -small, and B_3 -large graded unipotent Chevalley groups presented by the “weak” Steinberg relations are isomorphic to the groups presented by the “full” Steinberg relations.

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More specifically, we showed that the A_3 , B_3 -small, and B_3 -large graded unipotent Chevalley groups presented by the “weak” Steinberg relations are isomorphic to the groups presented by the “full” Steinberg relations.

Our proof strategy was to show that each full relation could be derived from the weak relations. We derived each relation by solving one or more group rewriting problems.

Why we proved it

Why we proved it

 Cornell University

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Mathematics > Group Theory

[Submitted on 8 Nov 2024]

Coboundary expansion inside Chevalley coset complex HDXs

Ryan O'Donnell, Noah G. Singer

Recent major results in property testing~\cite{BLM24,DDL24} and PCPs~\cite{BMV24} were unlocked by moving to high-dimensional expanders (HDXs) constructed from $\widetilde{C}_d$ -type buildings, rather than the long-known $\widetilde{A}_d$ -type ones. At the same time, these building quotient HDXs are not as easy to understand as the more elementary (and more symmetric/explicit) `\emph{coset complex}` HDXs constructed by Kaufman--Oppenheim~\cite{KO18} (of A_d -type) and O'Donnell--Pratt~\cite{OP22} (of B_d -, C_d -, D_d -type). Motivated by these considerations, we study the B_3 -type generalization of a recent work of Kaufman--Oppenheim~\cite{KO21}, which showed that the A_3 -type coset complex HDXs have good 1-coboundary expansion in their links, and thus yield 2-dimensional topological expanders.

The crux of Kaufman--Oppenheim's proof of 1-coboundary expansion was: (1)~identifying a group-theoretic result by Biss and Dasgupta~\cite{BD01} on small presentations for the A_3 -unipotent group over $\mathbb{F}_q$; (2)~``lifting" it to an analogous result for an A_3 -unipotent group over polynomial extensions $\mathbb{F}_q[x]$.

For our B_3 -type generalization, the analogue of~(1) appears to not hold. We manage to circumvent this with a significantly more involved strategy: (1)~getting a computer-assisted proof of vanishing 1-cohomology of B_3 -type unipotent groups over $\mathbb{F}_5$; (2)~developing significant new ``lifting" technology to deduce the required quantitative 1-cohomology results in B_3 -type unipotent groups over $\mathbb{F}_{5^k}[x]$.

Comments: 130 pages

Subjects: **Group Theory (math.GR)**; Discrete Mathematics (cs.DM)

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- To verify a result involving hundreds of calculations
 - Constants, negative signs, delicate computations
 - Noah: “What if I made a mistake somewhere?”
- To lay the groundwork for similar verifications
 - Lots of future work left!

Why work on this problem?



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- To construct higher-dimensional expanders

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Higher-dimensional expanders
(Specifically, topological expanders)

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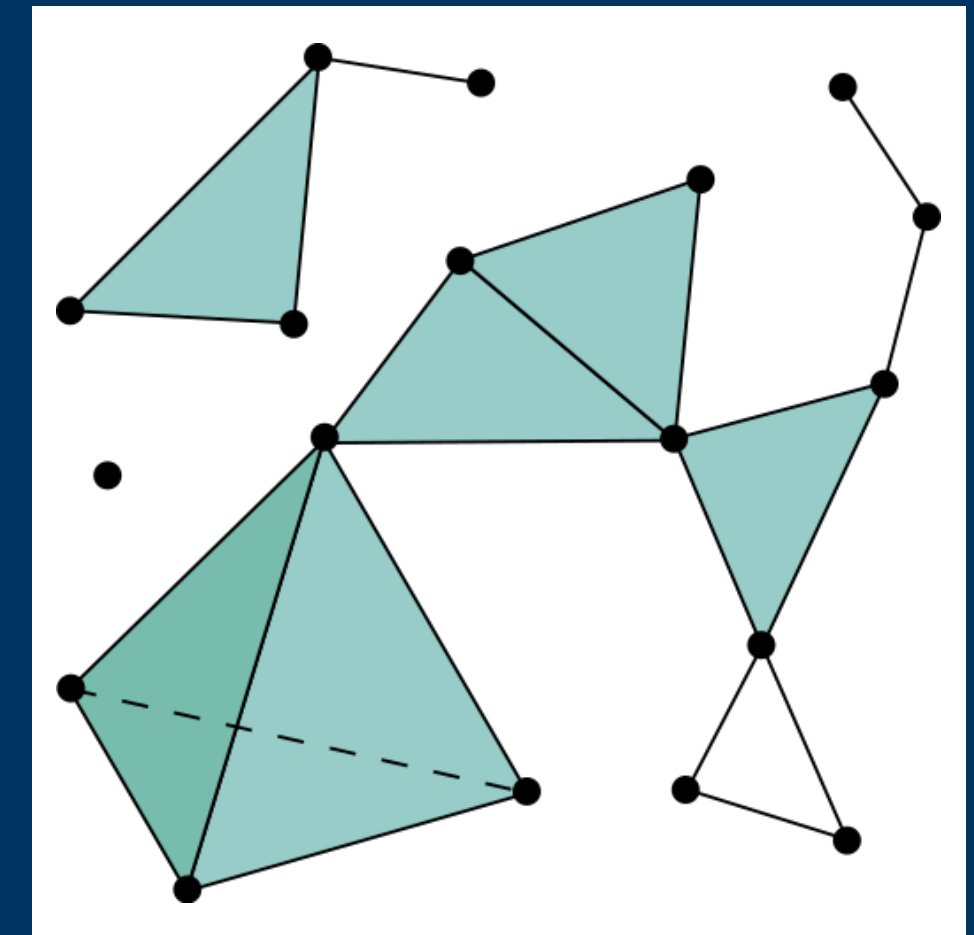
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Simplicial complexes



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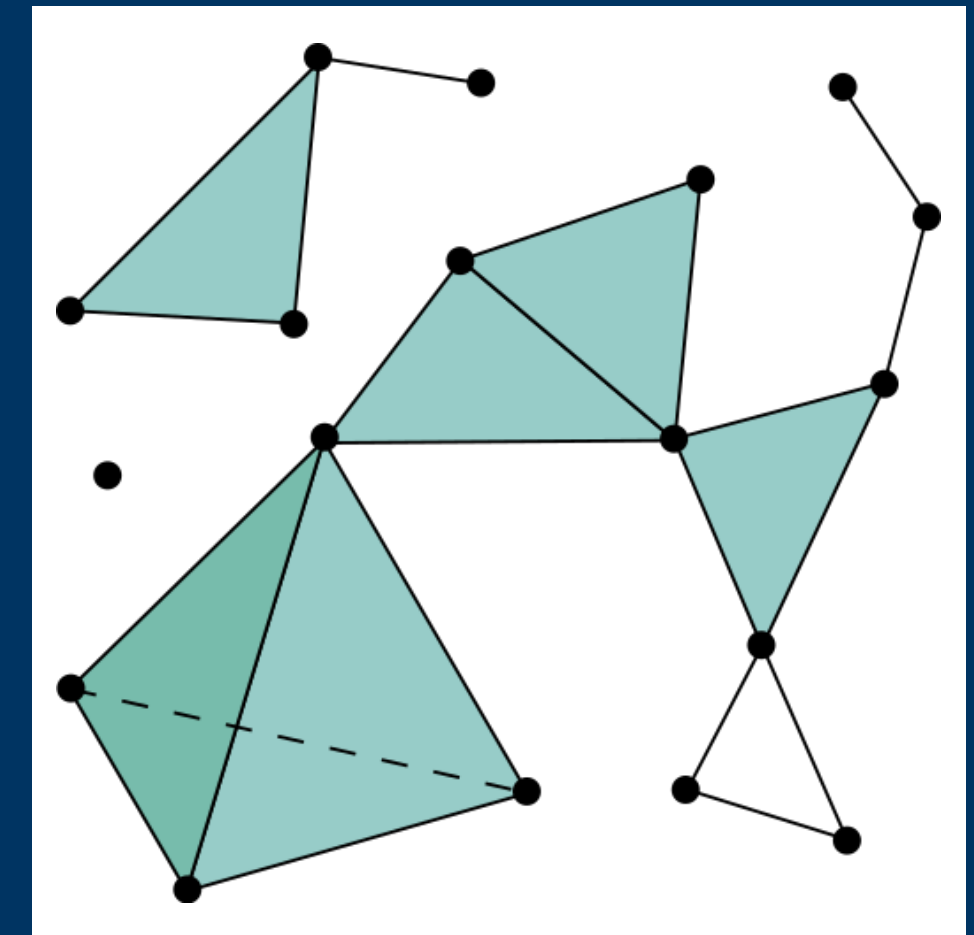
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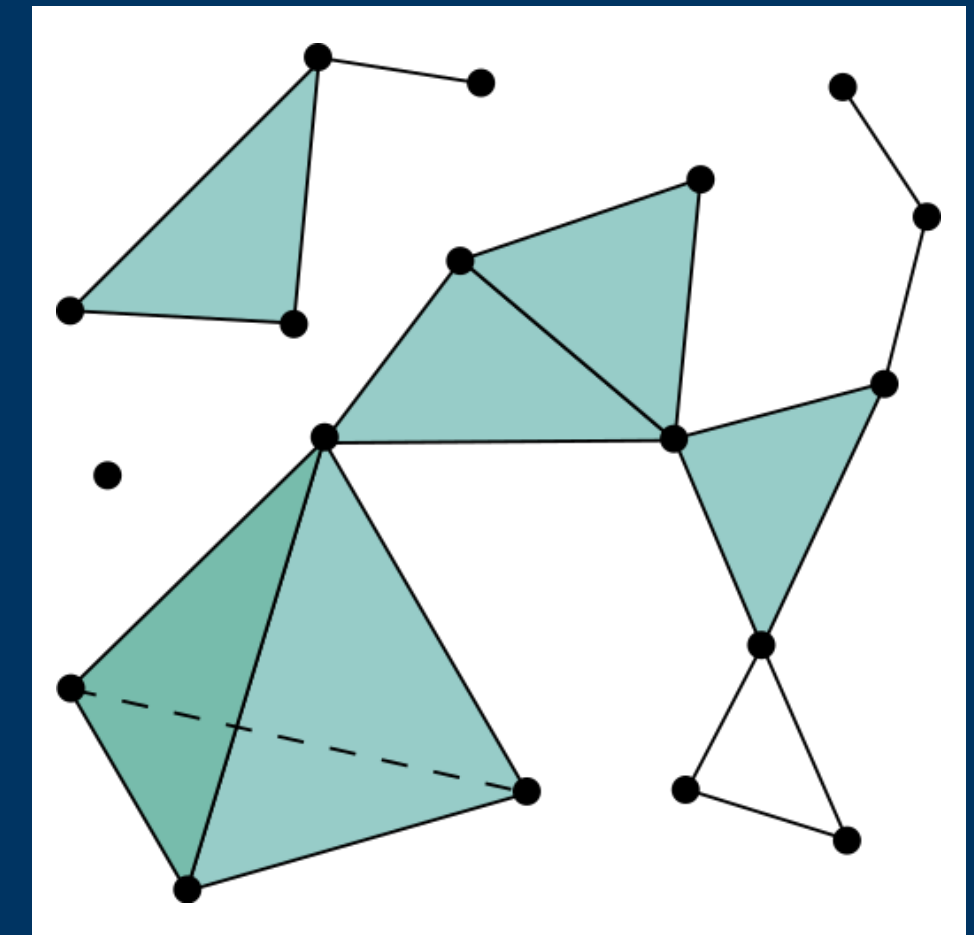
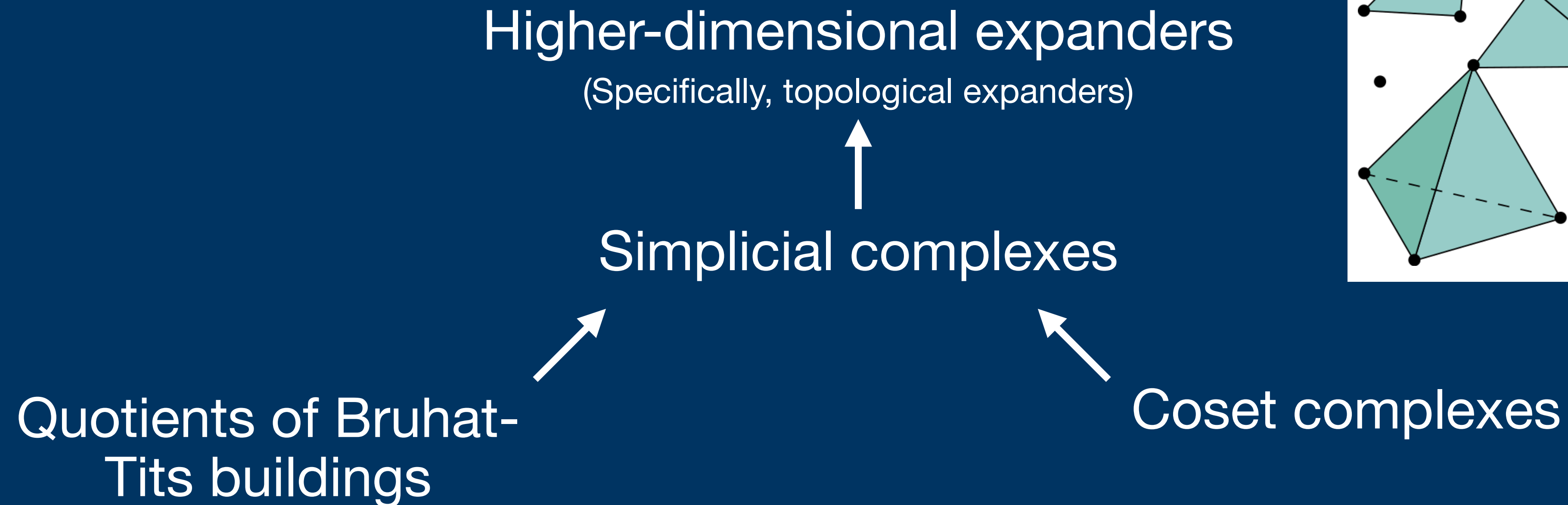
Quotients of Bruhat-Tits buildings



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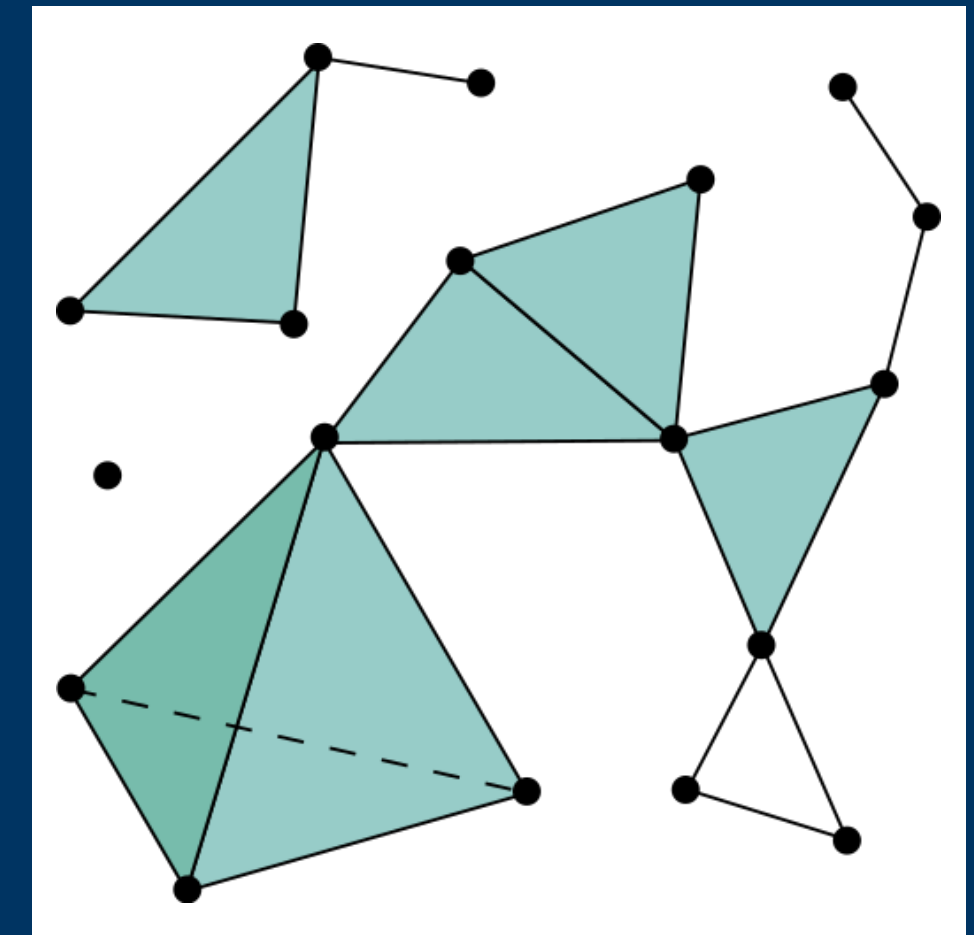
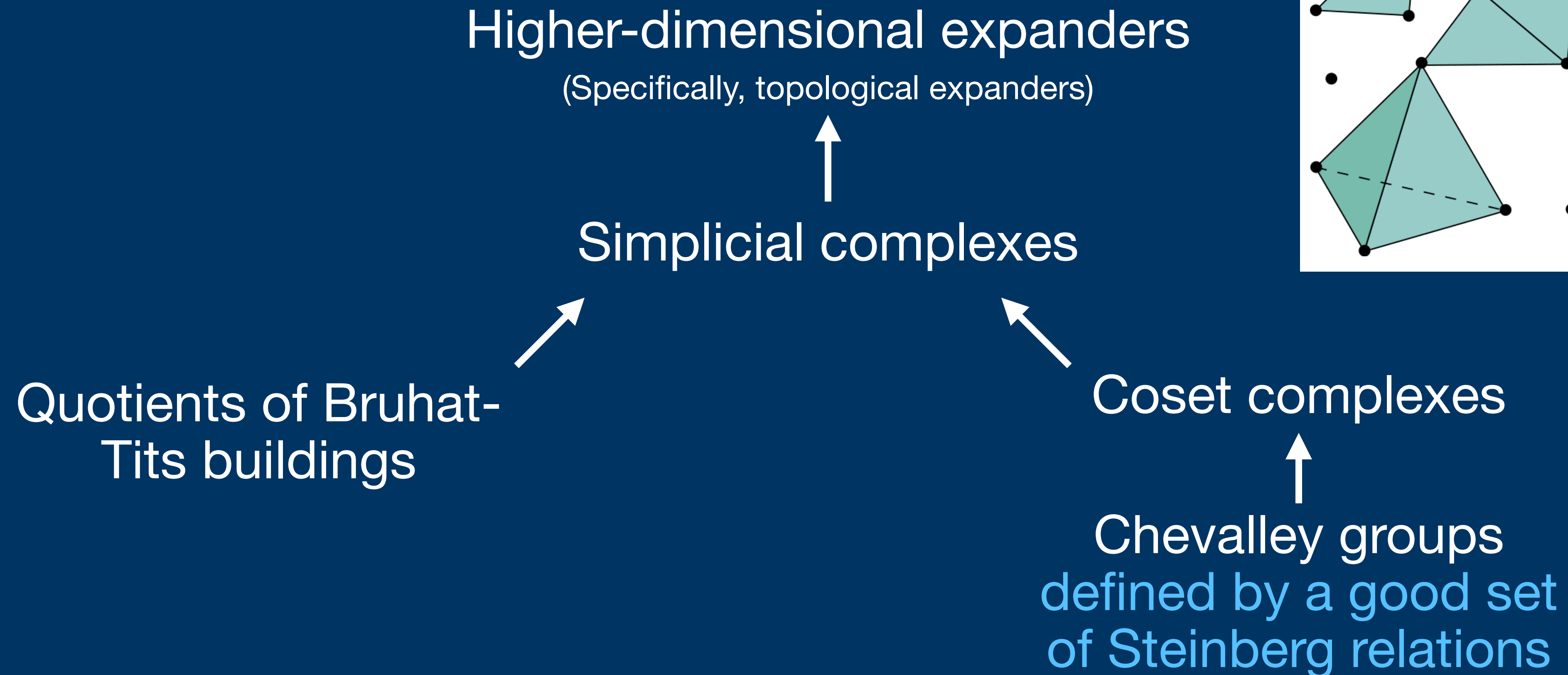
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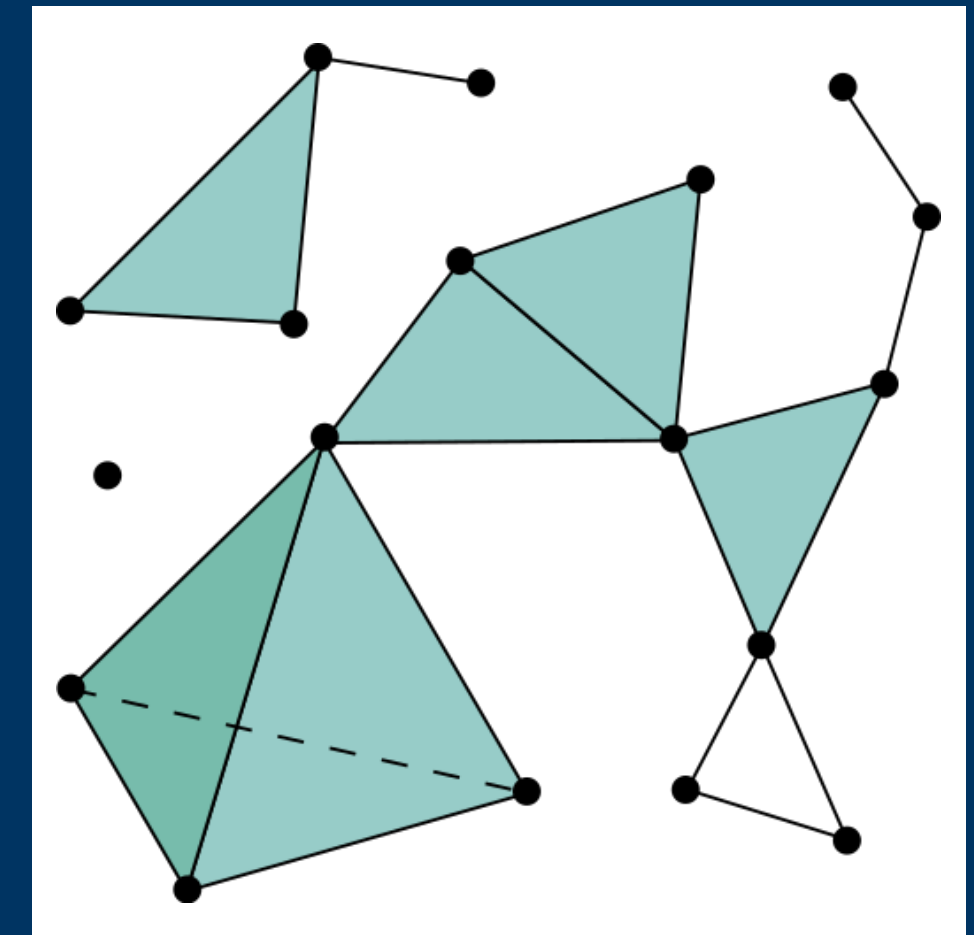
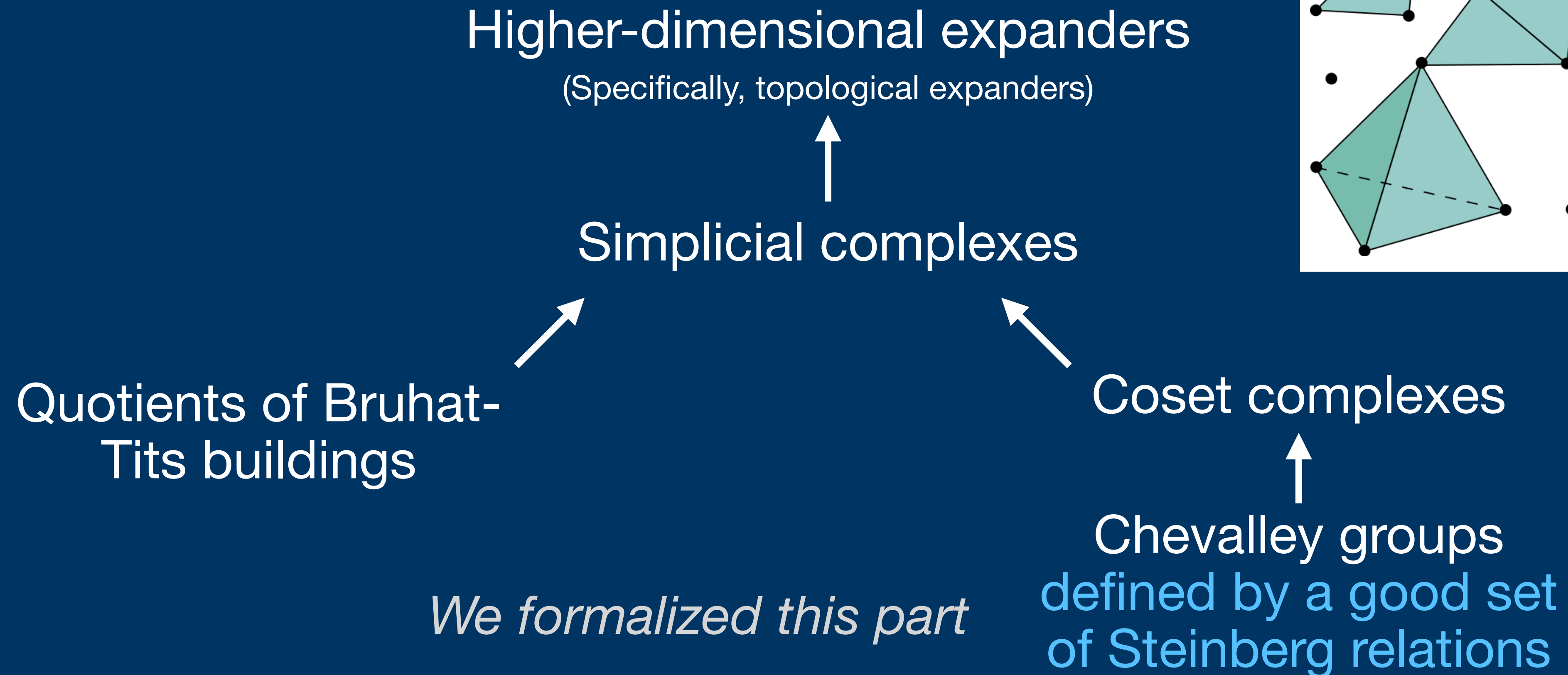
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 - Useful for (quantum) error correction, local property testing, and higher-dimensional geometry

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- To construct higher-dimensional expanders
 - Useful for (quantum) error correction, local property testing, and higher-dimensional geometry
- To do basic research in group theory

Chevalley groups



Chevalley groups



- Similar to Lie groups

Chevalley groups



- Similar to Lie groups
- Defined on square matrices containing field elements

Chevalley groups



- Similar to Lie groups
- Defined on square matrices containing field elements
- Express nice symmetries and geometric properties

Chevalley groups

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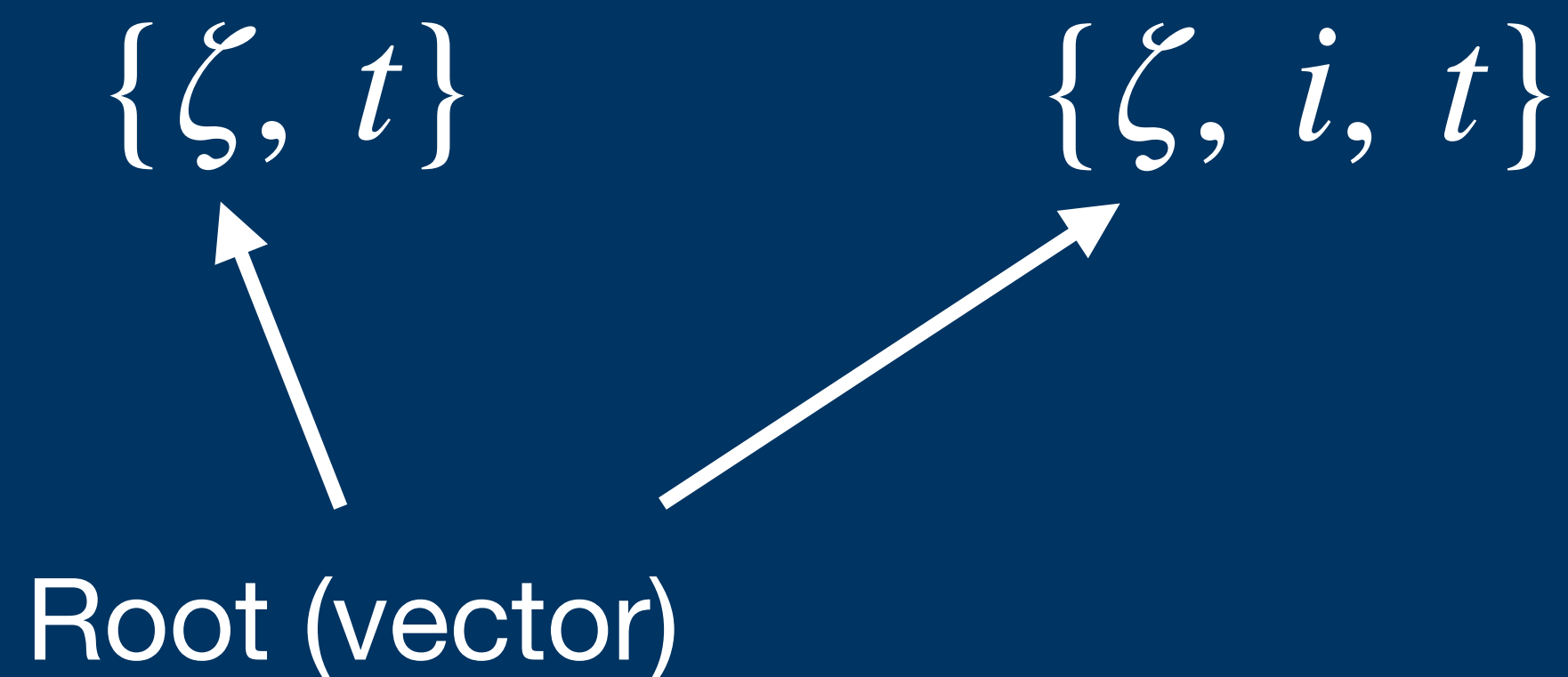
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- In our paper, these generators are either pairs or triples:

$$\{\zeta, t\}$$

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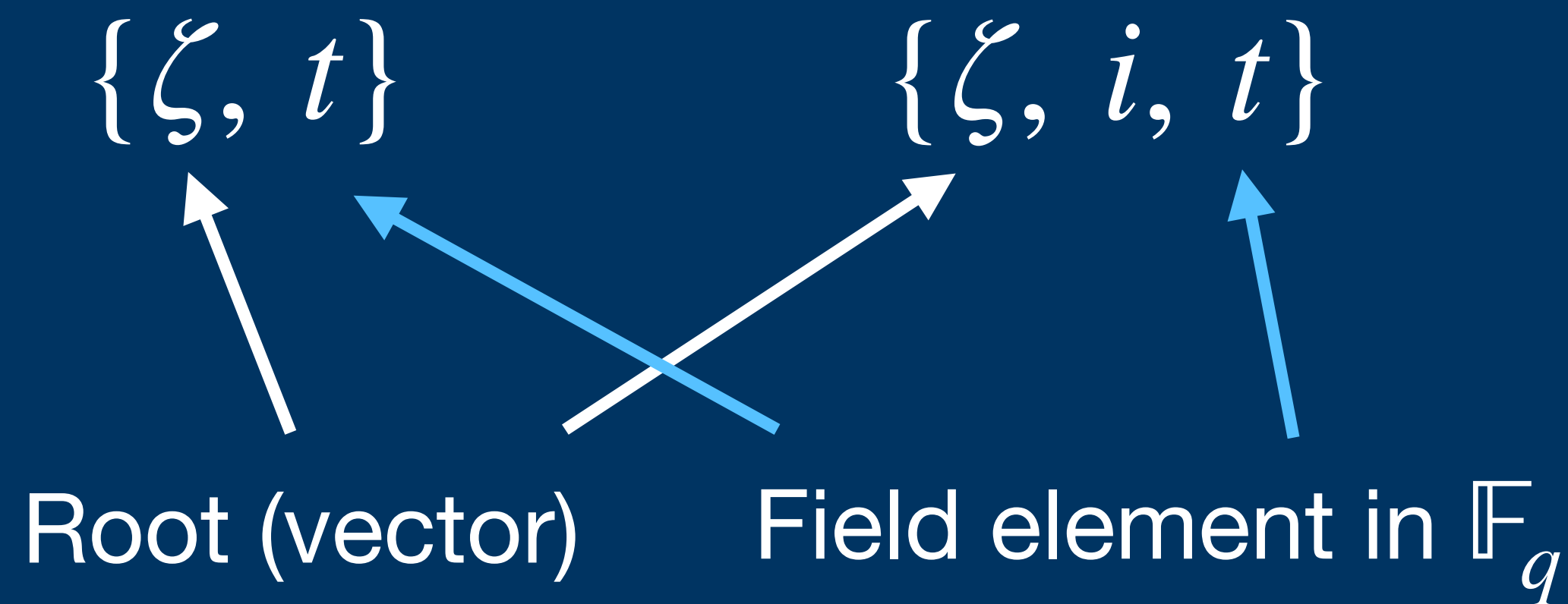
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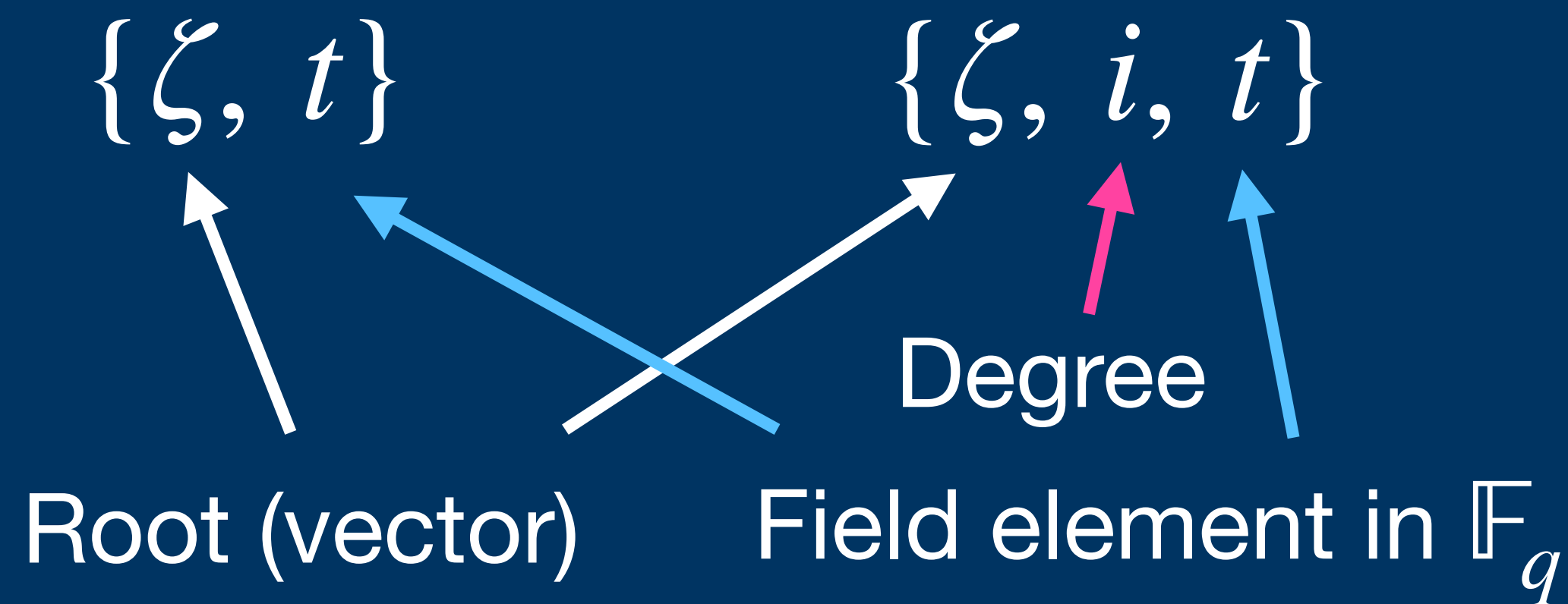
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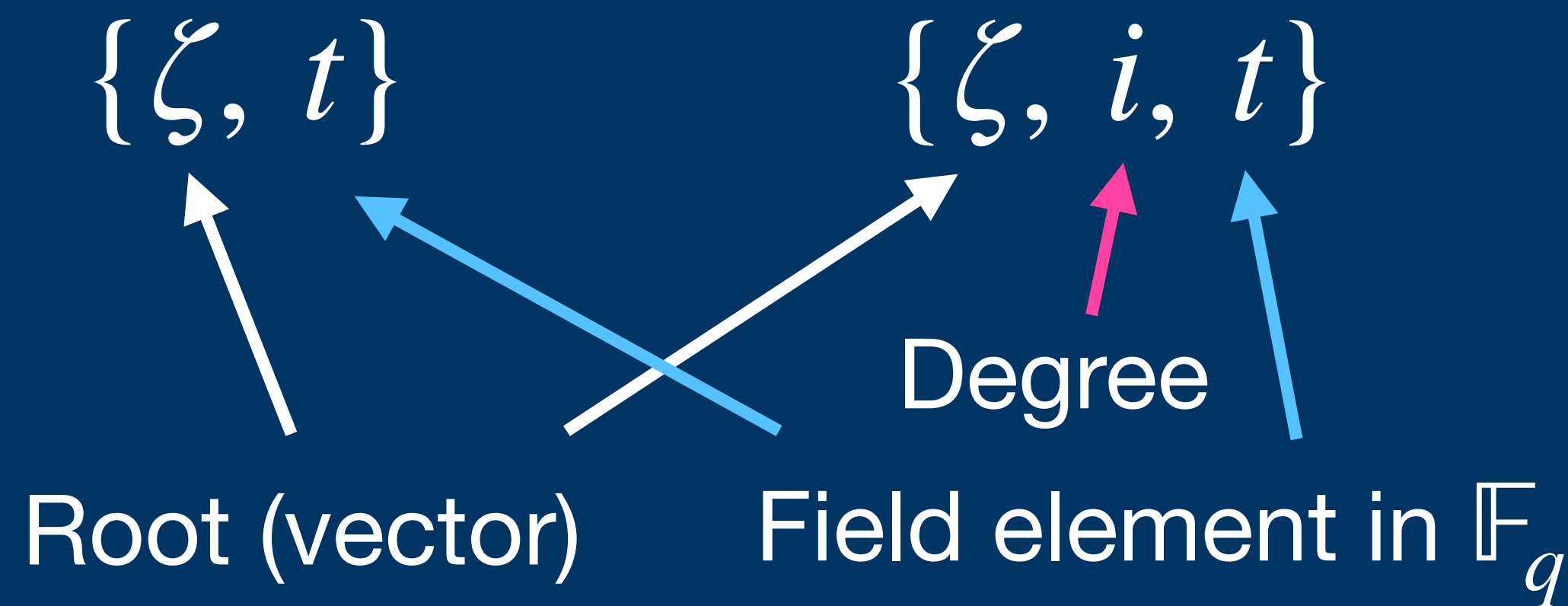
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```
structure GradedChevalleyGenerator
  (Φ : Type TΦ) [PositiveRootSystem Φ]
  (R : Type TR) [Ring R]
where
  ζ : Φ -- root
  i : ℕ -- degree
  hi : i ≤ height ζ
  t : R -- field coefficient
```

Chevalley groups

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- Commutator: For any pair of roots ζ and η , the commutator is a product of generators of the following form, for Chevalley constants a , b , and $C_{\zeta, \eta}^{a, b} \in [-3, 3]$:

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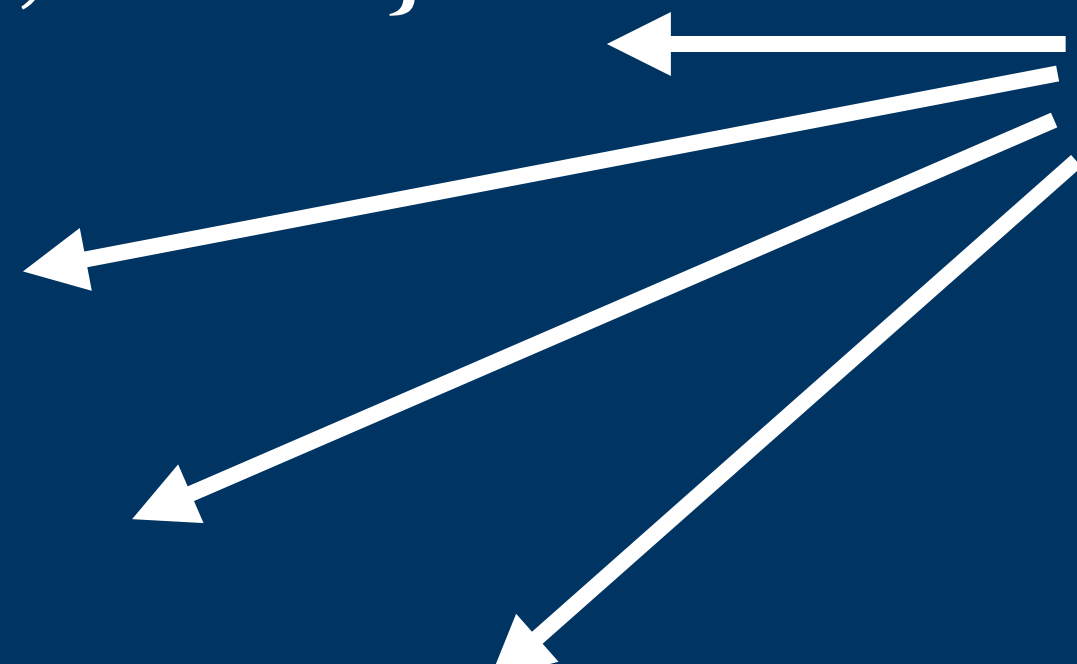
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These are (some of the) Steinberg relations



Presentations (not the academic kind)

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Equivalently, G is the free group on S modulo the normal closure of R in S

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An example calculation

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Assume the weak Steinberg relations. Then we want to show the full relation

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An example calculation

Assume the weak Steinberg relations. Then we want to show the full relation

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Proof. Decompose i arbitrarily into $i = i_1 + i_2$. Then ...

```
@[simp, chev_simps]
theorem lin_of_βψω : ∀ (i : N) (t u : F),
  {βψω, i, t} * {βψω, i, u} = {βψω, i, t + u} := by
  intro i hi t u
  rcases decompose i with ⟨i₁, i₂⟩      -- i = i₁ + i₂
  grw [ expr_βψω_as_β_ψω_β_ψω,          -- Line 1
        expr_βψω_ψω_as_ψω_βψω, expr_βψω_β_as_β_βψω, -- Line 2
        expr_βψω_ψω_as_ψω_βψω,          -- Line 3
        mul_one u, expr_βψω_as_β_ψω_β_ψω, -- Line 5
        mul_one (t + u), expr_βψω_as_β_ψω_β_ψω ] -- Line 6
```

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Proof. Decompose i arbitrarily into $i = i_1 + i_2$. Then ...

```
@[simp, chev_simps]
theorem lin_of_βψω : lin_of_root((weakB3SmallGraded F).project, βψω) :=
by
  intro i hi t u
  rcases decompose 1 2 i hi with ⟨i₁, i₂, rfl, hi₁, hi₂⟩
  rw [←mul_one t, expr_βψω_as_β_ψω_β_ψω Fchar hi₁ hi₂]
  grw [←expr_βψω_ψω_as_ψω_βψω Fchar, ←expr_βψω_β_as_β_βψω Fchar,
  ←expr_βψω_ψω_as_ψω_βψω Fchar]
  rw [←mul_one u]
  grw [expr_βψω_as_β_ψω_β_ψω Fchar hi₁ hi₂]
  rw [←mul_one (t + u)]
  grw [expr_βψω_as_β_ψω_β_ψω Fchar hi₁ hi₂]
  ring_nf
```

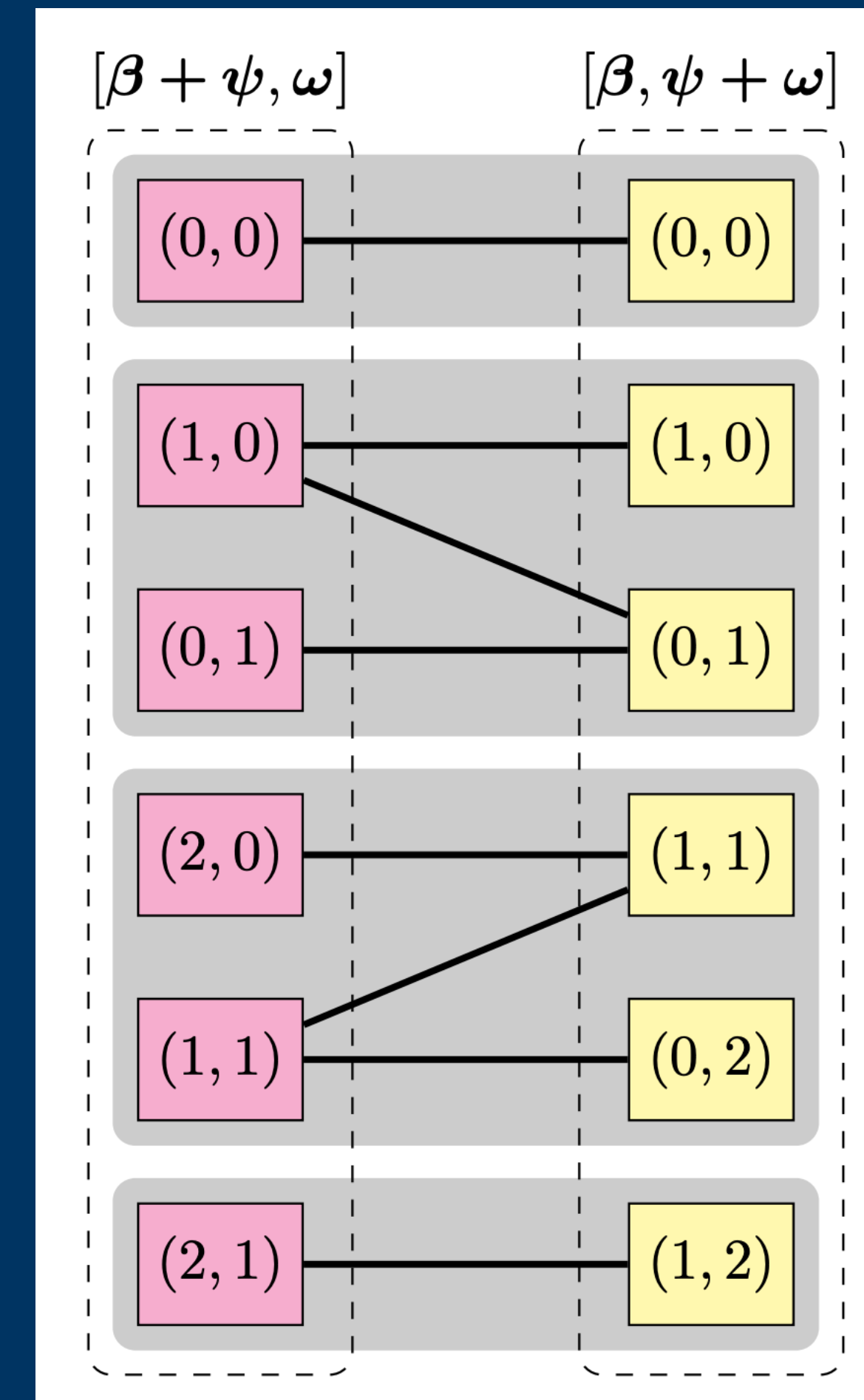
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Sometimes it's worse:

```
theorem expand_βψω_as_commutator_of_βψ_ω :  
  forall_ij_tu 2 1, ((βψω, i + j, 2 * t * u)) = [(βψ, i, t), (ω, j, u)] := by  
  intro i j hi hj t u  
  match i, j with  
  | 0, 0 => rw [expr_βψω_as_comm_of_βψ_ω_00 Fchar]  
  | 0, 1 => rw [expr_βψω_as_comm_of_βψ_ω_01 Fchar]  
  | 1, 0 => rw [expr_βψω_as_comm_of_βψ_ω_10 Fchar]  
  | 1, 1 => rw [expr_βψω_as_comm_of_βψ_ω_11 Fchar]  
  | 2, 0 => rw [expr_βψω_as_comm_of_βψ_ω_20 Fchar]  
  | 2, 1 => rw [expr_βψω_as_comm_of_βψ_ω_21 Fchar]
```

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Almost all of the proofs in our formalization look like that calculation.

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Sometimes it's better (i.e. we get proofs for free):

```
-- height 2 (reflection of height 1)
declare_B3Small_reflected_thm F b3small_valid  $\beta\psi\omega$   $\beta$   $\psi\omega$  const 1 heights 2 1 1 to 1 0 1
declare_B3Small_reflected_thm F b3small_valid  $\beta\psi\omega$   $\beta$   $\psi\omega$  const 1 heights 2 0 2 to 1 1 0
declare_B3Small_reflected_thm F b3small_valid  $\beta\psi\omega$   $\beta\psi$   $\omega$  const 2 heights 2 2 0 to 1 0 1
declare_B3Small_reflected_thm F b3small_valid  $\beta\psi\omega$   $\beta\psi$   $\omega$  const 2 heights 2 1 1 to 1 1 0
declare_B3Small_reflected_thm F b3small_valid  $\beta\psi\omega$   $\beta$   $\psi\omega$  const 1 heights 3 1 2 to 0 0 0
declare_B3Small_reflected_thm F b3small_valid  $\beta\psi\omega$   $\beta\psi$   $\omega$  const 2 heights 3 2 1 to 0 0 0
```

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inductive B3SmallPosRoot
```

```
| β | ψ | ω | βψ | ψω | β2ψ | βψω
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def height : B3SmallPosRoot → Nat
```

```
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*Some type isomorphisms
are better than others*

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```
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declare_B3Small_lin_id_inv_thms F ω
declare_B3Small_lin_id_inv_thms F βψ
declare_B3Small_lin_id_inv_thms F ψω
declare_B3Small_lin_id_inv_thms F β2ψ

declare_B3Small_trivial_span_of_root_pair_thms F β βψ
declare_B3Small_trivial_span_of_root_pair_thms F β β2ψ
declare_B3Small_trivial_span_of_root_pair_thms F ψ β2ψ
declare_B3Small_trivial_span_of_root_pair_thms F βψ β2ψ
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declare_B3Small_trivial_span_of_root_pair_thms F ω ψω

declare_B3Small_single_span_of_root_pair_thms F ψ ω ψω 2
declare_B3Small_single_span_of_root_pair_thms F ψ βψ β2ψ 2

/-! ### Mixed-degree theorem for specific roots -/

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Problem: need to declare the macros for each root system

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$$\{\beta, i, t\} \cdot \{\beta, i, -t\} \implies \{\beta, i, 0\} \implies 1$$

Other pain points

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Presentation relations vs. equations

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$$[\{\beta, i, t\}, \{\beta + \psi, j, u\}] = 1$$

$$\Leftrightarrow$$

$$\{\beta, i, t\} \cdot \{\beta + \psi, j, u\} \cdot \{\beta, i, -t\} \cdot \{\beta + \psi, j, -u\} = 1$$

$$\Leftrightarrow$$

$$\{\beta, i, t\} \cdot \{\beta + \psi, j, u\} = \{\beta + \psi, j, u\} \cdot \{\beta, i, t\}$$

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Group/field arithmetic often needed massaging to discharge with automation

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```
have aux1 : 2 * (u / (2 * t)) = u / t := by ring_nf; field_simp; group
have aux2 : u / (2 * t) * (u / (2 * t)) = (u * u) / (4 * (t * t)) := by
  ring_nf; simp only [inv_pow, mul_eq_mul_left_iff, inv_inj, mul_eq_zero, ne_eq,
    0fNat.ofNat_ne_zero, not_false_eq_true, pow_eq_zero_iff, inv_eq_zero];
  left
  rw [pow_two, mul_two, two_add_two_eq_four]
```

Thank you for your attention

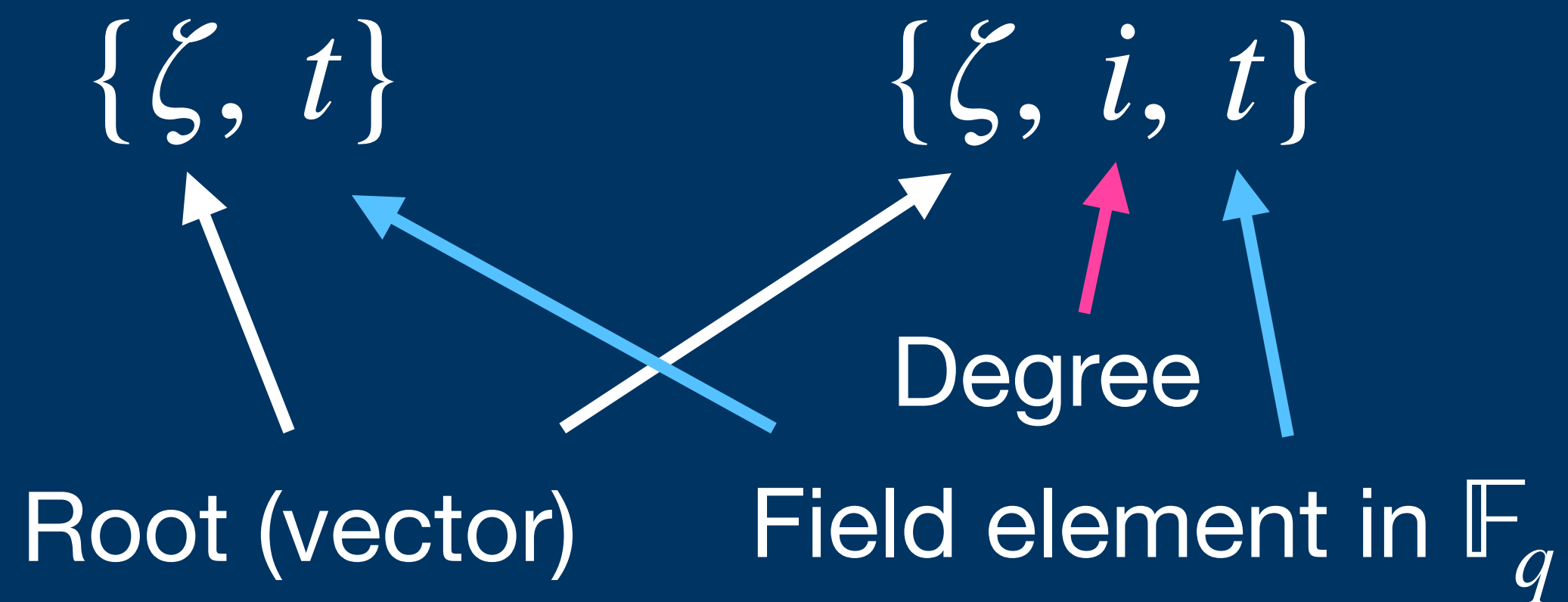
<https://github.com/singerng/steinberg-formalization>

$$\alpha \mid \beta \mid \alpha + \beta$$

theorem `lin_of_α`

$$a * (b * c) = (a * b) * c$$

lin_id_inv_thms



What we proved

