

Nondeterministic Asynchronous Dataflow in Isabelle/HOL

Rafael Castro G. Silva, Laouen Fernet and, Dmitriy Traytel

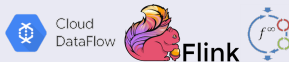
Department of Computer Science
University of Copenhagen

28/09/2025

Motivation

Context:

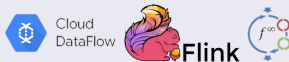
- Stream Processing: programs that compute with (possibly) unbounded sequences of data
- Frameworks: Google Cloud Dataflow, Apache Flink, and Timely Dataflow



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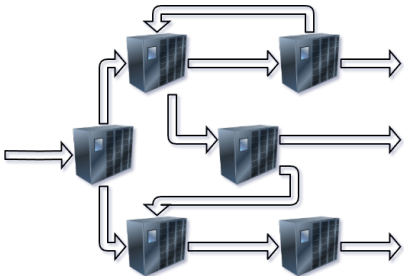
Our long term goal:

Mechanically Verify Timely Dataflow algorithms

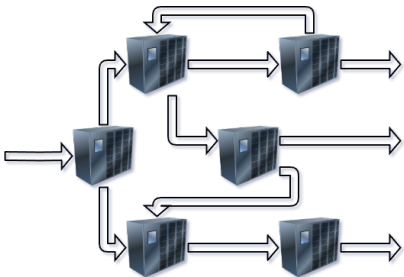
A Good Foundation

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- Nondeterministic Asynchronous Dataflow

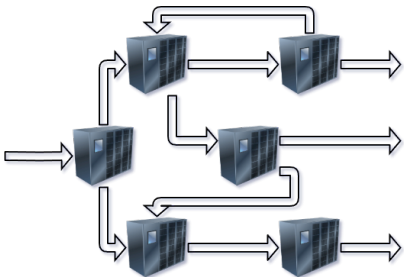


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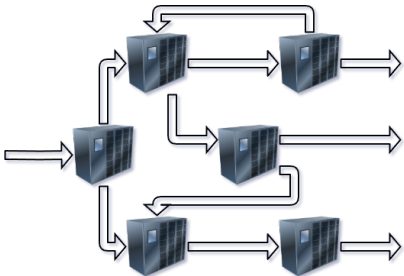
- Nondeterministic Asynchronous Dataflow
 - Dataflow: Directed graph of interconnected operators
 - Operators are (possibly) non-terminating processes that communicate with the network
 - Networks are unbounded FIFO queues

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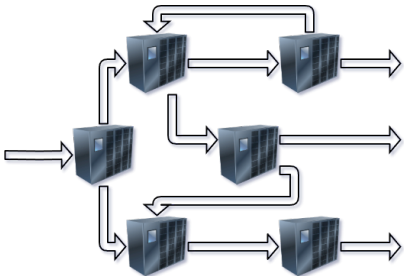
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Question

How do we know something is Nondeterministic Asynchronous Dataflow?

The Algebra for Nondeterministic Asynchronous Dataflow

- Bergstra et al. presents an algebra for Nondeterministic Asynchronous Dataflow
- Primitives:
sequential and parallel composition; feedback loop...
- 52 axioms
- A process calculus instance

Network Algebra for Asynchronous Dataflow*

J.A. Bergstra^{1,2,†} C.A. Middelburg^{2,3,§} Gh. Ştefănescu^{4,‡}

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⁴Institute of Mathematics of the Romanian Academy
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- An Isabelle/HOL instance of Nondeterministic Asynchronous Dataflow
 - Operators as a shallow embedding as codatatypes
 - 51 axioms proved

Operators as a Codatatype

Operators

Operators in Isabelle/HOL

codatatype ('i,'o,'d) *op* =

Read 'i ('d $\Rightarrow$ ('i,'o,'d) *op*) | Write (('i,'o,'d) *op*) 'o 'd |

Silent (('i,'o,'d) *op*) | Choice (('i,'o,'d) *op*) *cset*

(2, 1, nat) *op* shape



- Type parameters:
inputs/output ports; data
- Operator's actions
- Possibly infinite trees

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Uncommunicative operators

$\oslash = \text{Choice } \{\}_c$

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Uncommunicative operators

$\emptyset = \text{Choice } \{\}$ _c

$\odot = \text{Silent } \odot$

$\otimes = \text{Choice } \{\otimes\}$ _c

More examples

ex1 = Choice {Write ex1 1 42, $\emptyset$ }_c

ex2 = Choice {Write ex2 1 42, ex2}_c

ex3 = Choice {Write ex3 1 42, Silent ex3}_c

Operators Equivalences: Weak Bisimilarity

An ideal equivalence relation

- Only equate operators that **behave** the same
- Useful reasoning principle

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- Milner's approach!
The classic chapter on **weak bisimilarity**
- Based on labeled transition systems (LTS)
- **Weak**: silent computations are abstracted away

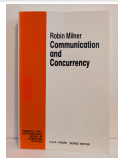


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Weak bisimilarity of operators

- Labels (actions): read, write, silent (τ)
- Weak bisimilarity of operators ($\approx$): LTS + weak simulation + greatest weak bisimulation
- $\approx$ has a useful coinduction principle
- $\emptyset \approx \odot$, $\odot \approx \otimes$, $\text{ex1} \approx \text{ex2}$, and $\text{ex2} \approx \text{ex3}$

Asynchronous Dataflow Operators

Auxiliary Definitions

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- Buffers: $'p \Rightarrow 'd$ list

Buffer functions

BHD $p \text{ buf} = \text{hd } (\text{buf } p)$

BTL $p \text{ buf} = \text{buf}(p := \text{tl } (\text{buf } p))$

BENQ $p \times \text{buf} = \text{buf}(p := \text{buf } p @ [x])$

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Mapping ports functions

$\text{map_op} :: ('i_1 \Rightarrow 'i_2) \Rightarrow ('o_1 \Rightarrow 'o_2) \Rightarrow ('i_1, 'o_1, 'd) \text{ op} \Rightarrow ('i_2, 'o_2, 'd) \text{ op}$

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$\curvearrowright :: ('a + 'b) + 'c \Rightarrow 'a + 'b + 'c$

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Equation of the identity operator

$\text{id_op} :: ('p \Rightarrow 'd \text{ list}) \Rightarrow ('p, 'p, 'd) \text{ op}$

$\text{id_op } \text{buf} = \text{Choice}$

$\cup_c$

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$$(((\lambda p. \text{Read } p (\lambda x. \text{id_op } (\text{BENQ } p \times buf)))) `c \ \mathcal{U}_c)$$
$$\cup_c$$

- $\mathcal{U}_c$ is the set of usable ports provide by its type

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$$((\lambda p. \text{Write } (\text{id_op } (\text{BTL } p \text{ } buf)) \text{ } p (\text{BHD } p \text{ } buf)) \text{`}_c \{p \in_c \mathcal{U}_c \mid buf \text{ } p \neq []\}))$$

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- $\mathcal{U}_c$ is the set of usable ports provide by its type
- Stream delayer: always can read, and it may eventually output

Identity

Equation of the identity operator

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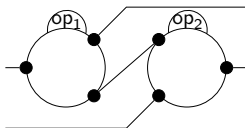
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Identity operator with an empty buffer

$$\mathcal{I} = \text{id_op } (\lambda_. [])$$

Composition: Preliminaries

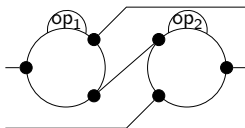
Composition operator type

$$\text{comp_op} :: ('o_1 \Rightarrow 'i_2 \text{ option}) \Rightarrow ('i_2 \Rightarrow 'd \text{ list}) \Rightarrow ('i_1, 'o_1, 'd) \text{ op} \Rightarrow ('i_2, 'o_2, 'd) \text{ op} \Rightarrow ('i_1 + 'i_2, 'o_1 + 'o_2, 'd) \text{ op}$$


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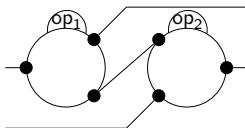
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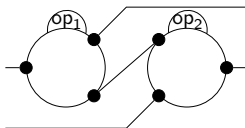
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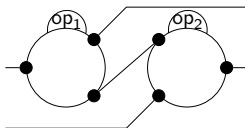
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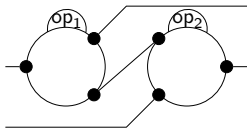
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Sequential and parallel composition

$$op_1 \parallel op_2 = \text{comp_op } (\lambda _ . \text{None}) (\lambda _ . []) op_1 op_2$$
$$op_1 \bullet op_2 = \text{map_op projl projr } (\text{comp_op } \text{Some } (\lambda _ . []) op_1 op_2)$$

Composition: Equation

Equation of the composition operator

$$\text{comp_op } \textit{wire} \textit{ buf } op_1 \textit{ op}_2 =$$

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$\text{comp_op } \textit{wire} \textit{ buf } op_1 \textit{ op}_2 = \text{Choice}$

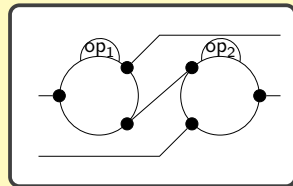
$\cup_c$

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$\text{comp_op wire buf } op_1 \ op_2 = \text{Choice } (((\lambda op. \underline{\text{case}} \ op \ \underline{\text{of}}$

$\text{'}_c \text{ choices } op_1) \cup_c$

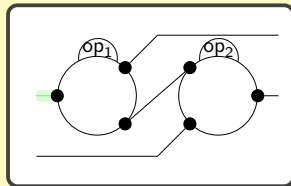


Composition: Equation

Equation of the composition operator

$$\text{comp_op wire buf } op_1 \text{ } op_2 = \text{Choice } (((\lambda op. \text{case } op \text{ of}$$
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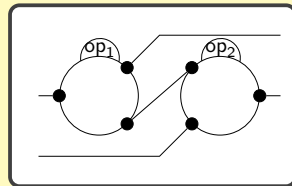


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 $| \text{Write } op \text{ } p \text{ } x \Rightarrow (\underline{\text{case}} \text{ wire } p \text{ of}$

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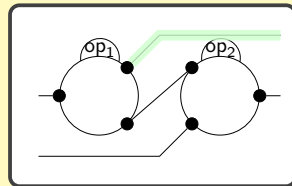


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Equation of the composition operator

$\text{comp_op wire buf } op_1 \ op_2 = \text{Choice } (((\lambda op. \underline{\text{case}} \ op \ \text{of}} \\ \text{Read } p \ f \Rightarrow \text{Read } (\text{Inl } p) \ (\lambda x. \text{comp_op wire buf } (f \ x) \ op_2) \\ | \text{Write } op \ p \ x \Rightarrow (\underline{\text{case}} \ \text{wire } p \ \text{of}} \\ \text{None} \Rightarrow \text{Write } (\text{comp_op wire buf } op \ op_2) (\text{Inl } p) \ x$

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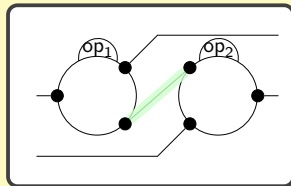


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 $| \text{Some } q \Rightarrow \text{Silent } (\text{comp_op wire } (\text{BENQ } q \ x \text{ buf}) op op_2)))$

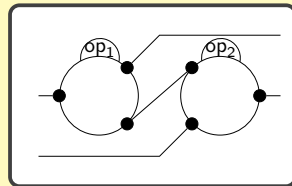
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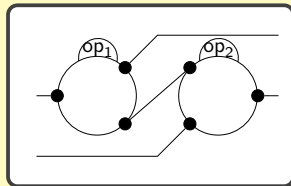
$\text{comp_op wire buf } op_1 op_2 = \text{Choice } (((\lambda op. \underline{\text{case}} op \text{ of}$
 $\text{Read } p f \Rightarrow \text{Read (Inl } p) (\lambda x. \text{comp_op wire buf } (f x) op_2)$
 $| \text{Write } op p x \Rightarrow (\underline{\text{case}} wire p \text{ of}$
 $\text{None} \Rightarrow \text{Write (comp_op wire buf } op op_2) (\text{Inl } p) x$
 $| \text{Some } q \Rightarrow \text{Silent (comp_op wire (BENQ } q x \text{ buf) } op op_2)))$
 $| \text{Silent } op \Rightarrow \text{Silent (comp_op wire buf } op op_2))$
 $\backslash_c \text{ choices } op_1) \cup_c$



Composition: Equation

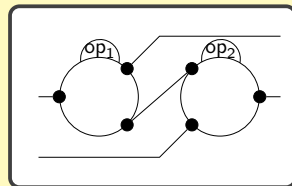
Equation of the composition operator

$\text{comp_op wire buf } op_1 op_2 = \text{Choice } (((\lambda op. \underline{\text{case}} op \text{ of}$
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 $| \text{ Silent } op \Rightarrow \text{Silent (comp_op wire buf } op op_2))$
 $\text{`}_c \text{ choices } op_1) \cup_c$



Composition: Equation

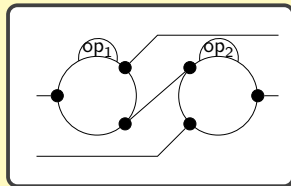
Equation of the composition operator

$$\begin{aligned}
 \text{comp_op wire buf } op_1 \text{ } op_2 &= \text{Choice } (((\lambda op. \text{case } op \text{ of} \\
 &\quad \text{Read } p \text{ } f \Rightarrow \text{Read } (\text{Inl } p) (\lambda x. \text{comp_op wire buf } (f \text{ } x) \text{ } op_2) \\
 &\quad | \text{Write } op \text{ } p \text{ } x \Rightarrow (\text{case wire } p \text{ of} \\
 &\quad \quad \text{None} \Rightarrow \text{Write } (\text{comp_op wire buf } op \text{ } op_2) (\text{Inl } p) \text{ } x \\
 &\quad \quad | \text{Some } q \Rightarrow \text{Silent } (\text{comp_op wire } (\text{BENQ } q \text{ } x \text{ } buf) \text{ } op \text{ } op_2))) \\
 &\quad | \text{Silent } op \Rightarrow \text{Silent } (\text{comp_op wire buf } op \text{ } op_2)) \\
 &\quad \text{`}_c \text{ choices } op_1) \cup_c \\
 &\quad ((\lambda op. \text{case } op \text{ of}
 \end{aligned}$$


$$\text{`}_c \text{ sound_reads wire buf (choices } op_2)))$$

Composition: Equation

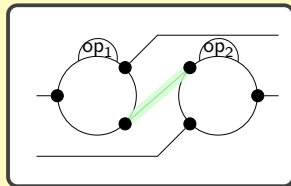
Equation of the composition operator

$$\begin{aligned}
 \text{comp_op wire buf } op_1 op_2 = & \text{Choice } (((\lambda op. \text{case } op \text{ of} \\
 & \text{Read } p f \Rightarrow \text{Read } (\text{Inl } p) (\lambda x. \text{comp_op wire buf } (f x) op_2) \\
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 & | \text{Silent } op \Rightarrow \text{Silent } (\text{comp_op wire buf } op op_2)) \\
 & \text{`}_c \text{ choices } op_1) \cup_c \\
 & ((\lambda op. \text{case } op \text{ of} \\
 & \quad \text{Read } p f \Rightarrow \text{if } p \in \text{ran wire}
 \end{aligned}$$


$$\text{`}_c \text{ sound_reads wire buf } (\text{choices } op_2)))$$

Composition: Equation

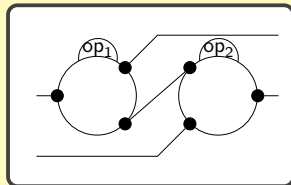
Equation of the composition operator

$$\begin{aligned}
 \text{comp_op wire buf } op_1 op_2 = & \text{Choice } (((\lambda op. \text{case } op \text{ of} \\
 & \text{Read } p \ f \Rightarrow \text{Read } (\text{Inl } p) (\lambda x. \text{comp_op wire buf } (f \ x) op_2) \\
 & | \text{Write } op \ p \ x \Rightarrow (\text{case wire } p \text{ of} \\
 & \quad \text{None} \Rightarrow \text{Write } (\text{comp_op wire buf } op \ op_2) (\text{Inl } p) \ x \\
 & \quad | \text{Some } q \Rightarrow \text{Silent } (\text{comp_op wire } (\text{BENQ } q \ x \ \text{buf}) op \ op_2))) \\
 & | \text{Silent } op \Rightarrow \text{Silent } (\text{comp_op wire buf } op \ op_2)) \\
 \text{`}_c \text{ choices } op_1) \cup_c & \\
 ((\lambda op. \text{case } op \text{ of} & \\
 \text{Read } p \ f \Rightarrow \text{if } p \in \text{ran wire} & \\
 \text{then Silent } (\text{comp_op wire } (\text{BTL } p \ \text{buf}) op_1 & (f \ (\text{BHD } p \ \text{buf}))))
 \end{aligned}$$


$$\text{`}_c \text{ sound_reads wire buf } (\text{choices } op_2)))$$

Composition: Equation

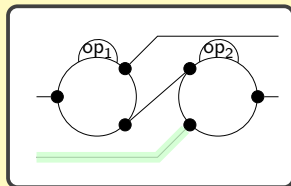
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 \text{comp_op wire buf } op_1 op_2 = & \text{Choice } (((\lambda op. \underline{\text{case}} \text{ } op \text{ of} \\
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 \end{aligned}$$


$$\text{`}_c \text{ sound_reads wire buf } (\text{choices } op_2)))$$

Composition: Equation

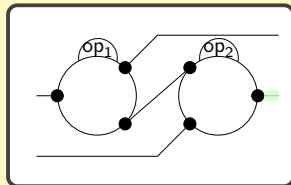
Equation of the composition operator

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 \text{comp_op wire buf } op_1 op_2 = & \text{Choice } (((\lambda op. \text{case } op \text{ of} \\
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 \text{`}_c \text{ choices } op_1) \cup_c & \\
 ((\lambda op. \text{case } op \text{ of} & \\
 \text{Read } p \ f \Rightarrow \text{if } p \in \text{ran wire} & \\
 \text{then Silent } (\text{comp_op wire } (\text{BTL } p \ \text{buf}) \ op_1 \ (f \ (\text{BHD } p \ \text{buf}))) & \\
 \text{else Read } (\text{Inr } p) \ (\lambda x. \text{comp_op wire buf } op_1 \ (f \ x)) &
 \end{aligned}$$


$$\text{`}_c \text{ sound_reads wire buf } (\text{choices } op_2)))$$

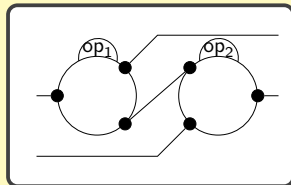
Composition: Equation

Equation of the composition operator

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 \text{comp_op wire buf } op_1 op_2 = & \text{Choice } (((\lambda op. \text{case } op \text{ of} \\
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 \text{`}_c \text{ choices } op_1) \cup_c & \\
 ((\lambda op. \text{case } op \text{ of} & \\
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 & \quad \text{then Silent } (\text{comp_op wire } (\text{BTL } p \ \text{buf}) \ op_1 \ (f \ (\text{BHD } p \ \text{buf}))) \\
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 & \text{`}_c \text{ choices } op_1) \cup_c \\
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 & \quad \text{Read } p \ f \Rightarrow \text{if } p \in \text{ran wire} \\
 & \quad \quad \underline{\text{then}} \text{Silent (comp_op wire (BTL } p \text{ buf) } op_1 (f (\text{BHD } p \text{ buf}))) \\
 & \quad \quad \underline{\text{else}} \text{Read (Inr } p) (\lambda x. \text{comp_op wire buf } op_1 (f \ x)) \\
 & | \text{Write } op \ p \ x \Rightarrow \text{Write (comp_op wire buf } op_1 \text{ op) (Inr } p) \ x \\
 & | \text{Silent } op \Rightarrow \text{Silent (comp_op wire buf } op_1 \text{ op)) `}_c \text{ sound_reads wire buf (choices } op_2)))
 \end{aligned}$$


Asynchronous Dataflow Operators

- Identity: $\mathcal{I}$
- Parallel composition: $op_1 \parallel op_2$
- Sequential composition: $op_1 \bullet op_2$

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Transpose operator type

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Dummy source operator type

$i :: ('i, 'o, 'd) op$

Sink operator type

$! :: ('i, 'o, 'd) op$

Split operator type

$\Lambda :: ('n, 'n + 'n, 'd) op$

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Transpose operator type

$\mathcal{X} :: ('n + 'm, 'm + 'n, 'd) op$

Dummy source operator type

$i :: ('i, 'o, 'd) op$

Sink operator type

$! :: ('i, 'o, 'd) op$

Split operator type

$\Lambda :: ('n, 'n + 'n, 'd) op$

Merge operator type

$\mathcal{V} :: ('n + 'n, 'n, 'd) op$

Asynchronous Dataflow Operators

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$\uparrow :: ('i + 'm, 'o + 'm, 'd) op \Rightarrow ('i, 'o, 'd) op$

Transpose operator type

$\mathcal{X} :: ('n + 'm, 'm + 'n, 'd) op$

Dummy source operator type

$i :: ('i, 'o, 'd) op$

Sink operator type

$! :: ('i, 'o, 'd) op$

Split operator type

$\Lambda :: ('n, 'n + 'n, 'd) op$

Merge operator type

$\mathcal{V} :: ('n + 'n, 'n, 'd) op$

Copy operator type

$\mathcal{C} :: ('n, 'n + 'n, 'd) op$

Asynchronous Dataflow Operators

- Identity: $\mathcal{I}$
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Feedback operator type ($op \uparrow$)

$\uparrow :: ('i + 'm, 'o + 'm, 'd) op \Rightarrow ('i, 'o, 'd) op$

Transpose operator type

$\mathcal{X} :: ('n + 'm, 'm + 'n, 'd) op$

Dummy source operator type

$i :: ('i, 'o, 'd) op$

Sink operator type

$! :: ('i, 'o, 'd) op$

Split operator type

$\Lambda :: ('n, 'n + 'n, 'd) op$

Merge operator type

$\mathcal{V} :: ('n + 'n, 'n, 'd) op$

Copy operator type

$\mathcal{C} :: ('n, 'n + 'n, 'd) op$

Equality test operator type

$\mathcal{Q} :: ('n + 'n, 'n, 'd \text{ option}) op$

Asynchronous Dataflow Properties

Basic Network Algebra Properties

- B1: $op_1 \parallel (op_2 \parallel op_3) \approx \text{map_op } \curvearrowright \curvearrowright (op_1 \parallel op_2) \parallel op_3$
B2_1: $op \parallel (\mathcal{I} :: (0, 0, 'd) op) \approx \text{map_op } \text{Inl } \text{Inl } op$
B2_2: $(\mathcal{I} :: (0, 0, 'd) op) \parallel op \approx \text{map_op } \text{Inr } \text{Inr } op$
B3: $(op_1 \bullet op_2) \bullet op_3 \approx op_1 \bullet (op_2 \bullet op_3)$
B4_1: $op \bullet \mathcal{I} \approx op$ B4_2: $\mathcal{I} \bullet op \approx op$
B5: $(op_1 \parallel op_2) \bullet (op_3 \parallel op_4) \approx (op_1 \bullet op_3) \parallel (op_2 \bullet op_4)$
B6: $\mathcal{I} \parallel \mathcal{I} \approx \mathcal{I}$ B7: $\mathcal{X} \bullet \mathcal{X} \approx \mathcal{I}$
B8: $(\mathcal{X} :: ('i + 0, 0 + 'i, 'd) op) \approx \text{map_op } \text{id } (\text{case_sum } \text{Inr } \text{Inl}) \mathcal{I}$
B9: $\mathcal{X} \approx \text{map_op } \curvearrowright \curvearrowright (\mathcal{X} \parallel \mathcal{I}) \bullet \text{map_op } \text{id } \curvearrowright (\mathcal{I} \parallel \mathcal{X})$
B10: $(op_1 \parallel op_2) \bullet \mathcal{X} \approx \mathcal{X} \bullet (op_2 \parallel op_1)$
F1: $\mathcal{I} \uparrow \approx (\mathcal{I} :: (0, 0, 'd) op)$ F2: $\mathcal{X} \uparrow \approx \mathcal{I}$
R1: $op_2 \bullet (op_1 \uparrow) \approx ((op_2 \parallel \mathcal{I}) \bullet op_1) \uparrow$
R2: $(op_1 \uparrow) \bullet op_2 \approx (op_1 \bullet (op_2 \parallel \mathcal{I})) \uparrow$
R3: $op_1 \parallel (op_2 \uparrow) \approx (\text{map_op } \curvearrowright \curvearrowright (op_1 \parallel op_2)) \uparrow$
R4: $(op_1 \bullet (\mathcal{I} \parallel op_2)) \uparrow \approx ((\mathcal{I} \parallel op_2) \bullet op_1) \uparrow$
R5: $\text{map_op } \text{Inl } \text{Inl } ((op :: ('i + 0, 'o + 0, 'd) op) \uparrow) \approx op$
R6: $(op \uparrow) \uparrow \approx (\text{map_op } \curvearrowright \curvearrowright op) \uparrow$

Basic Network Algebra Properties

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- B9: $\mathcal{X} \approx \text{map_op } \curvearrowright \curvearrowright (\mathcal{X} \parallel \mathcal{I}) \bullet \text{map_op } \text{id } \curvearrowright (\mathcal{I} \parallel \mathcal{X})$
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- R3: $op_1 \parallel (op_2 \uparrow) \approx (\text{map_op } \curvearrowleft \curvearrowleft (op_1 \parallel op_2)) \uparrow$
- R4: $(op_1 \bullet (\mathcal{I} \parallel op_2)) \uparrow \approx ((\mathcal{I} \parallel op_2) \bullet op_1) \uparrow$
- R5: $\text{map_op } \text{Inl } \text{Inl } ((op :: ('i + 0, 'o + 0, 'd) op) \uparrow) \approx op$
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R1: $op_2 \bullet (op_1 \uparrow) \approx ((op_2 \parallel \mathcal{I}) \bullet op_1) \uparrow$
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R5: $\text{map_op } \text{Inl } \text{Inl } ((op :: ('i + 0, 'o + 0, 'd) op) \uparrow) \approx op$
R6: $(op \uparrow) \uparrow \approx (\text{map_op } \curvearrowright \curvearrowright op) \uparrow$

Basic Network Algebra Properties

- B1: $op_1 \parallel (op_2 \parallel op_3) \approx \text{map_op } \curvearrowright \curvearrowright (op_1 \parallel op_2) \parallel op_3$
B2_1: $op \parallel (\mathcal{I} :: (0, 0, 'd) op) \approx \text{map_op } \text{Inl } \text{Inl } op$
B2_2: $(\mathcal{I} :: (0, 0, 'd) op) \parallel op \approx \text{map_op } \text{Inr } \text{Inr } op$
B3: $(op_1 \bullet op_2) \bullet op_3 \approx op_1 \bullet (op_2 \bullet op_3)$
B4_1: $op \bullet \mathcal{I} \approx op$ B4_2: $\mathcal{I} \bullet op \approx op$
B5: $(op_1 \parallel op_2) \bullet (op_3 \parallel op_4) \approx (op_1 \bullet op_3) \parallel (op_2 \bullet op_4)$
B6: $\mathcal{I} \parallel \mathcal{I} \approx \mathcal{I}$ B7: $\mathcal{X} \bullet \mathcal{X} \approx \mathcal{I}$
B8: $(\mathcal{X} :: ('i + 0, 0 + 'i, 'd) op) \approx \text{map_op } \text{id } (\text{case_sum } \text{Inr } \text{Inl}) \mathcal{I}$
B9: $\mathcal{X} \approx \text{map_op } \curvearrowright \curvearrowright (\mathcal{X} \parallel \mathcal{I}) \bullet \text{map_op } \text{id } \curvearrowright (\mathcal{I} \parallel \mathcal{X})$
B10: $(op_1 \parallel op_2) \bullet \mathcal{X} \approx \mathcal{X} \bullet (op_2 \parallel op_1)$
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A2: $\mathcal{X} \bullet \bigotimes \approx \bigotimes$ for $\bigotimes \in \{Q, V\}$ A6: $\bigotimes \bullet \mathcal{X} \approx \bigotimes$ for $\bigotimes \in \{C, \Lambda\}$
A3_Q: $((i :: (0, 'a, 'd) \text{ op}) \parallel I) \bullet Q \approx (! :: (0 + 'a, 0, 'd) \text{ op}) \bullet i$
A3_V: $((i :: (0, 'a, 'd) \text{ op}) \parallel I) \bullet V \approx \text{map_op lnr id } I$
A4: $\bigotimes \bullet ! \approx ! \parallel !$ for $\bigotimes \in \{Q, V\}$ A8: $i \bullet \bigotimes \approx i \parallel i$ for $\bigotimes \in \{C, \Lambda\}$
A5: $C \bullet (C \parallel I) \approx \text{map_op id } \curvearrowright (C \bullet (I \parallel C))$
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A10: $Q \bullet C \approx (C \parallel C) \bullet (\text{map_op } \curvearrowright \curvearrowright (\text{map_op } \curvearrowright \curvearrowright (I \parallel \mathcal{X}) \parallel I)) \bullet (Q \parallel Q)$
A11: $C \bullet Q \approx I$ A12: $i \approx (I :: (0, 0, 'd) \text{ op})$
A16: $! \approx (I :: (0, 0, 'd) \text{ op})$ A13: $i \approx i \parallel i$ A17: $! \approx ! \parallel !$
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A15: $\bigotimes \approx \text{map_op } \curvearrowright \curvearrowright (\text{map_op } \curvearrowright \curvearrowright (I \parallel \mathcal{X}) \parallel I) \bullet (\bigotimes \parallel \bigotimes)$ for $\bigotimes \in \{Q, V\}$
A18: $\text{map_op lnr id } (\bigotimes :: (0, 0 + 0, 'd) \text{ op}) \approx I$ for $\bigotimes \in \{C, \Lambda\}$
A19: $\bigotimes \approx (\bigotimes \parallel \bigotimes) \bullet \text{map_op } \curvearrowright \curvearrowright (\text{map_op } \curvearrowright \curvearrowright (I \parallel \mathcal{X}) \parallel I)$ for $\bigotimes \in \{C, \Lambda\}$
F3: $(\text{map_op id lnr } \bigotimes) \uparrow \approx !$ for $\bigotimes \in \{Q, V\}$
F4: $(\text{map_op lnr id } \bigotimes) \uparrow \approx i$ for $\bigotimes \in \{C, \Lambda\}$
F5: $((I \parallel C) \bullet \text{map_op } \curvearrowright \curvearrowright (\mathcal{X} \parallel I) \bullet (I \parallel Q)) \uparrow \approx ! \bullet i$

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$$A6: \bigotimes \bullet \mathcal{X} \approx \bigotimes \quad \text{for } \bigotimes \in \{C, \Lambda\}$$

$$A3_Q: ((i :: (0, 'a, 'd) \text{ op}) \parallel I) \bullet Q \approx (! :: (0 + 'a, 0, 'd) \text{ op}) \bullet i$$

$$A3_V: ((i :: (0, 'a, 'd) \text{ op}) \parallel I) \bullet V \approx \text{map_op } \text{Inr id } I$$

$$A4: \bigotimes \bullet ! \approx ! \parallel ! \quad \text{for } \bigotimes \in \{Q, V\}$$

$$A8: i \bullet \bigotimes \approx i \parallel i \quad \text{for } \bigotimes \in \{C, \Lambda\}$$

$$A5: C \bullet (C \parallel I) \approx \text{map_op id } \curvearrowright (C \bullet (I \parallel C))$$

$$A7: C \bullet (! \parallel I) \approx \text{map_op id } \text{Inr } I$$

$$A9: i \bullet ! \approx I$$

$$A10: Q \bullet C \approx (C \parallel C) \bullet (\text{map_op } \curvearrowright \curvearrowright (\text{map_op } \curvearrowleft \curvearrowleft (I \parallel \mathcal{X}) \parallel I)) \bullet (Q \parallel Q)$$

$$A11: C \bullet Q \approx I$$

$$A12: i \approx (I :: (0, 0, 'd) \text{ op})$$

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$$A13: i \approx i \parallel i$$

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$$A14: \text{map_op id } \text{Inl } (\bigotimes :: (0 + 0, 0, 'd) \text{ op}) \approx I \quad \text{for } \bigotimes \in \{Q, V\}$$

$$A15: \bigotimes \approx \text{map_op } \curvearrowright \curvearrowright (\text{map_op } \curvearrowleft \curvearrowleft (I \parallel \mathcal{X}) \parallel I) \bullet (\bigotimes \parallel \bigotimes) \quad \text{for } \bigotimes \in \{Q, V\}$$

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Properties of Equality Test, Merge, Copy, Split, Source and Sink operators

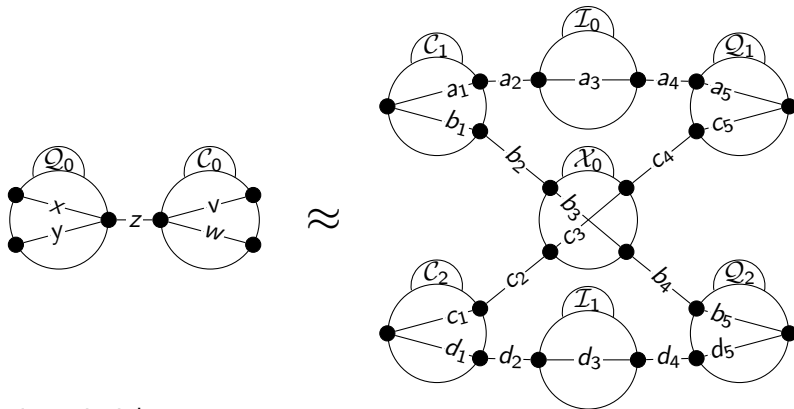
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- $\approx$ coinduction principle
- Buffers generalization
- Buffers invariant (details in the paper and artifact)

Related Work

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- Vélus project
 - Verified Lustre compilation
 - Synchronous dataflow language

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- Around 28 000 lines of definitions and proofs
- Future work:
 - Brock–Ackermann Time anomaly
 - Formalize Timely Dataflow infrastructure (finished)
 - Verify Timely Dataflow algorithms (on-going)

Questions, comments and suggestions