PART 1/2

EMDASH

A DEPENDENTLY TYPED LOGICAL FRAMEWORK FOR COMPUTATIONAL SYNTHETIC CATEGORY THEORY

Based on the paper by Author

INFORMAL MATHEMATICS

- Fluid & Intuitive
- Structural Reasoning



FORMAL SYSTEMS

- Rigid & Explicit
- Foundational Encoding

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CAN WE BRIDGE THIS GAP?

A DIFFERENT VISION: FUNCTORIAL PROGRAMMING

"The core entities of a mathematical domain—such as categories and functors—are not merely encoded but are themselves the primitive building blocks of the formal language."

Inspired by Kosta Dosen [1, 2]

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Proofs of structural integrity should be **computations**, not separate objects.

A novel dependently typed logical framework that implements these functorial principles.

• **Synthetic:** Categorical notions (`Cat`, `Obj`, `Functor`) are primitives.

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- **Specification-Driven:** Implemented in TypeScript, formally specified in Lambdapi.
- **Practical:** The formal engine for **hotdocX**, an Al-assisted formalization platform.

CORE ARCHITECTURE

TYPE THEORY CORE

- λΠ-Calculus Modulo Theory
- HOAS for binders
- Bidirectional Type Checker
- Unification & Hole Solving

SYNTHETIC PRIMITIVES

- Cat`, `Obj`, `Hom`
- `Functor`, `Transf`
- Yoneda (`hom_cov`)
- User-defined rewrite rules

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The system is designed to be extensible, allowing users to add their own definitions, rewrite rules, and unification <u>hints</u>.

THE KILLER FEATURE: FUNCTORIAL ELABORATION

THE TRADITIONAL WAY: DATA + PROOF

In Coq/Agda/Lean, you define a functor by providing:

- 1. **Data:** An object map F_0 and a morphism map F_1 .
- 2. **Proof:** A separate proof term `F_preserves_comp` demonstrating that $F_1(g \circ f) = F_1(g) \circ F_1(f)$.

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The structure and its laws are separate entities.

THE EMDASH WAY: DEFINITIONAL CHECK

In Emdash, defining a functor uses a single primitive:

MkFunctorTerm(C, D, fmap0, fmap1)

The elaboration engine itself **definitionally verifies** the coherence laws during type checking.

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Coherence is a check, not a proof.

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5. The Verdict:

- **Success:** Elaboration succeeds. The term is accepted.
- X **Failure:** Throw a `CoherenceError` with the non-equal normal forms.

BLUEPRINT FOR SYNTHESIS: THE LAMBDAPI SPECIFICATION

The implementation is not ad-hoc; it's guided by a formal, executable specification.

```
constant symbol Cat : TYPE;
injective symbol Obj : Cat → TYPE;
injective symbol Hom : Π [A : Cat] (X: Obj A) (Y: Obj A), TYPE;

// ...

constant symbol Functor_cat : Π(A : Cat), Π(B : Cat), Cat;

// Morphisms in the functor category ARE natural transformations
rule @Hom (Functor_cat _ _) $F $G $G$ Transf $F $G;

// Functoriality is a computational rule
rule compose_morph (@fapp1 _ _ $F _ _ $a) (@fapp1 _ _ $F _ _ $a')
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This declarative style makes coherence laws part of the system's core computational behavior.

INTERACTIVE THEOREM PROVING

Emdash isn't just for defining; it's for proving.

A proof-in-progress is a term with unsolved goals, called **Holes** (`?h0`).

```
Goal ?g0: \vdash \Pi (n : Nat). Nat
```

Users construct proofs by refining these holes using tactics.

EXAMPLE PROOF SESSION

Goal: Prove `id: Nat -> Nat`

```
// 1. State Goal
defineGlobal("id_nat_proof", Pi("n", Expl, Nat, _ => Nat), Hole("?g0"))
// > Goal ?g0: ⊢ Π (n : Nat). Nat

// 2. Introduce hypothesis
intro(Var("id_nat_proof"), "?g0", "n")
// > Hole ?g0 solved with `\lambda n. ?g1`
// > New Goal ?g1: n : Nat ⊢ Nat

// 3. Solve goal with exact term
exact(Var("id_nat_proof"), "?g1", Var("n"))
// > Hole ?g1 solved with `n`

// 4. Proof Complete!
// Final term: \(\lambda\) (n : Nat). n
```

VALIDATION & TESTING

Emdash is a working TypeScript implementation, verified by a comprehensive test suite.

- Inductive types (Nat, List, Vec) & dependent functions.
- **V** Higher-order unification & pattern matching.
- **Crucially:** Tests confirm that `MkFunctorTerm` accepts valid functors and throws `CoherenceError` for invalid ones.
- **V** Full interactive proof sessions are tested from start to finish.

VALIDATION & TESTING

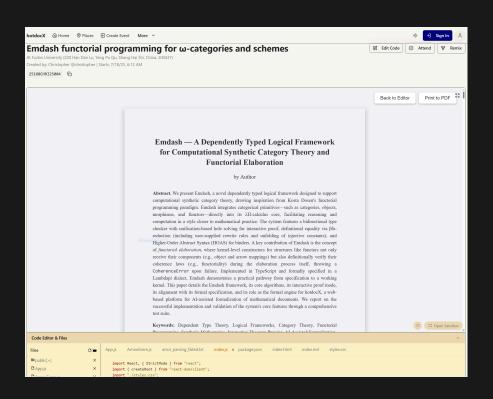
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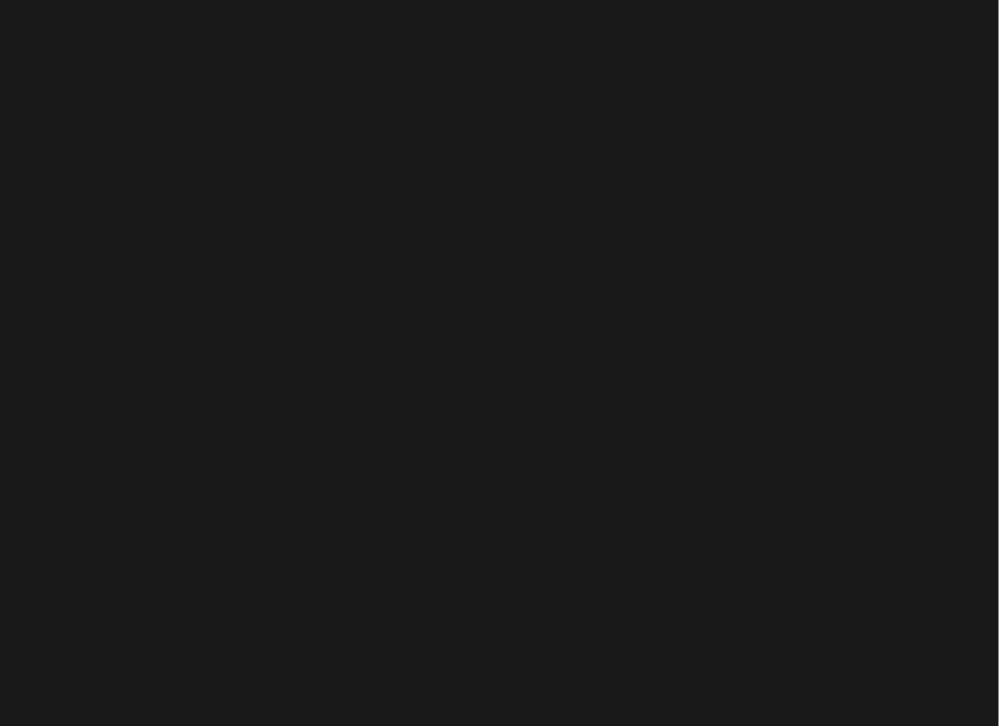
The system behaves as specified.

THE GRAND VISION: HOTDOCX

Emdash is the formal engine for **hotdocX**, a web platform for AI-assisted formalization.

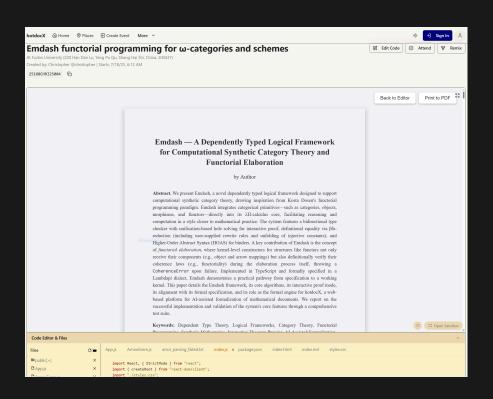


Goal: Transform mathematical documents into executable, verifiable, and interactive formal content.



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Goal: Transform mathematical documents into executable, verifiable, and interactive formal content.

Emdash's interactive proof mode and clear error reporting are designed for human-AI collaboration.

CONCLUSION

Emdash demonstrates a practical pathway for "functorial programming":

- A **synthetic** approach to category theory in a dependently typed framework.
- A novel **functorial elaboration** mechanism that makes coherence a definitional check.
- A robust implementation guided by a formal specification.
- A practical tool with interactive proving, serving as the kernel for the hotdocX platform.

FUTURE WORK

DEEPEN THE SYNTHESIS

- ω-Categories via `Total_cat`
- Limits, Adjunctions
- Univalence

ENHANCE THE FRAMEWORK

- Universe Hierarchy
- Performance Tuning
- Richer Tactic Language

HOTDOCX & AI

- Al-driven proof suggestions
- Document parsing (LaTeX)
- Visualization tools

THANK YOU

QUESTIONS?

Project Repository: github.com/hotdocx/emdash

hotdocX Platform: hotdocx.github.io

Continue: ./slides-part-2-hotdocx.html

PART 2/2

HOTDOCX & JSCOQ

AN INTERACTIVE, AI-AUGMENTED, AND MONETIZABLE PLATFORM FOR COQ

Coq is a cornerstone of formal methods, but sharing its rich, interactive proofs is hard.

CURRENT STATE

- Static PDFs (Flattened)
- Code Repositories



THE COMMUNITY'S DESIRE

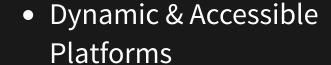
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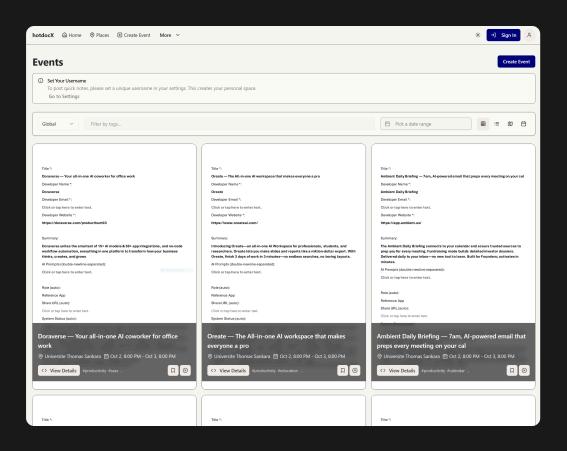
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THE CORE PROBLEM: A MISSING PLATFORM THAT IS BOTH INTERACTIVE AND SUSTAINABLE.

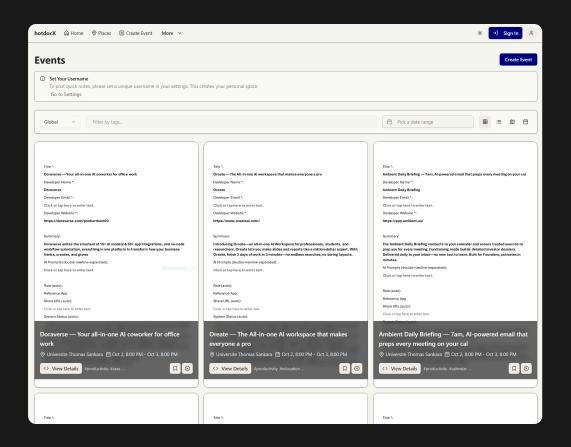
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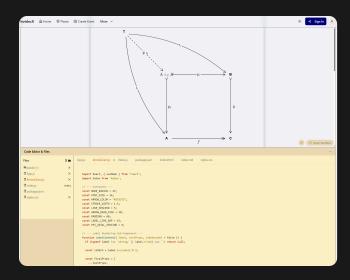


We're moving from "papers-with-code" to "papers-as-apps".

CORE CONCEPT: AI TEMPLATES

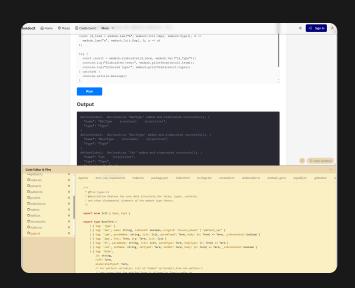
A user provides a document (PDF, LaTeX, etc.) and a prompt. Our AI pipeline uses a template to generate a runnable application.

Arrowgram Template



Renders commutative diagrams from natural language.

Emdash Template

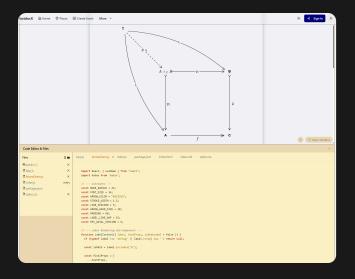


Provides a live environment for our logical framework.

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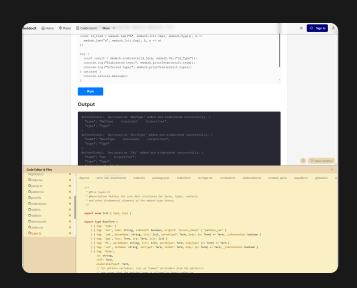
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The platform is flexible and context-aware. Now, let's add Coq to the mix.

MAKING COQ A FIRST-CLASS CITIZEN

We created a 'jsCoq' template for the hotdocX ecosystem.

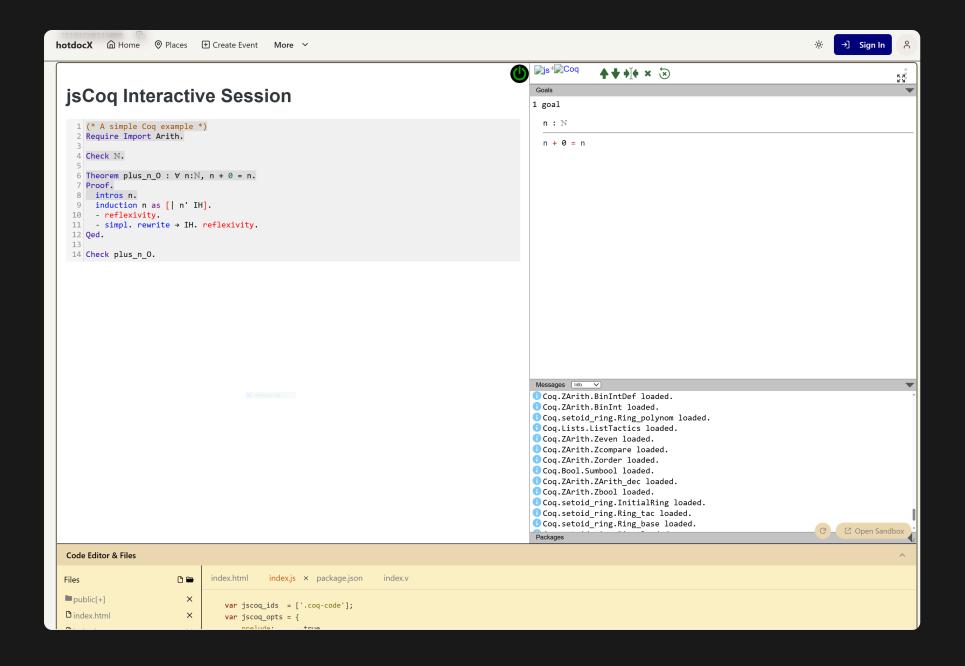
Any Coq script can become a first-class, interactive, and shareable hotdocX event.

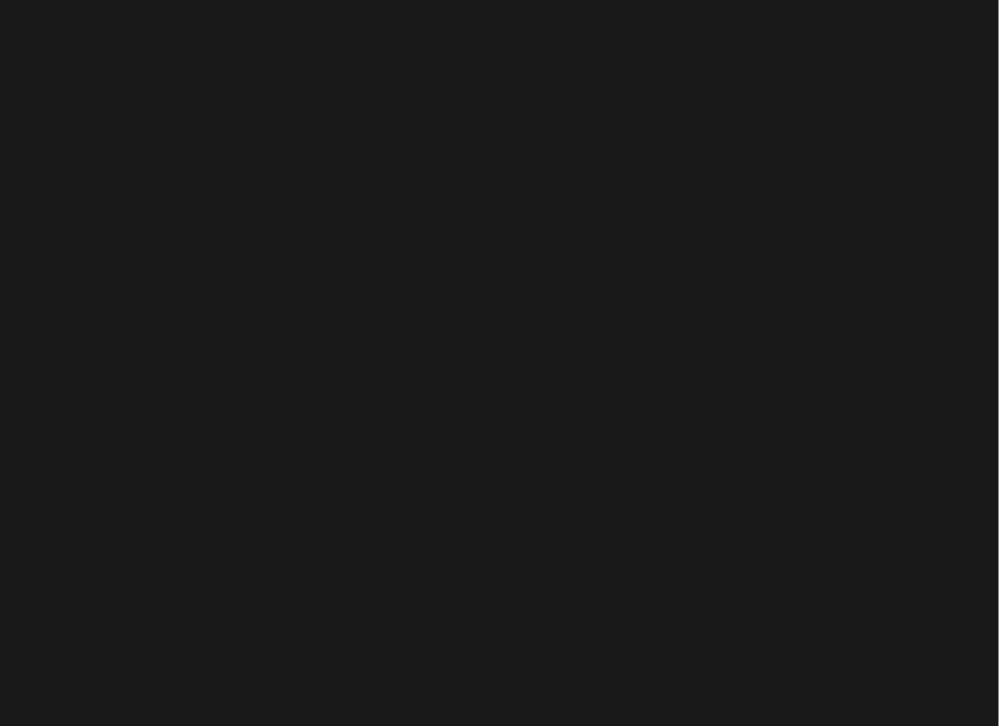
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LIVE DEMO 1: INTERACTIVE COQ SESSION





UNLOCKING POWERFUL NEW WORKFLOWS

EDUCATION

- Create tutorials from lectures + Coq files.
- Students solve exercises inbrowser.

RESEARCH

- Publish an arXiv paper and a companion hotdocX event.
- Readers can experiment with lemmas directly.

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- Bridge the gap between paper and formalization.

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The co-location of `jsCoq` and our native framework, `Emdash`, opens a unique research opportunity.

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EDUCATION	RESEARCH	AI INTEGRATION

A "hackable" kernel for teaching CIC without a complex build setup.

Rapidly prototype new tactics and language features in TypeScript.

A native JS kernel is easier for AI agents to interact with than a compiled binary.

We have a working platform, a clear vision, and a sustainable model (dynamic fee percentage on transactions, applied to new authors referred by white-labelled multi-tenant publishers-hosted instances of hotdocx; i.e. `professor1.github.io/hotdocx` with `publisher2.github.io/hotdocx`).

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Your contribution, whether individual or institutional, directly funds development, infrastructure, and community growth.

GITHUB.COM/SPONSORS/HOTDOCX

Sponsoring provides premium and VIP tiers on the hotdocX platform.

THANK YOU

QUESTIONS & DISCUSSION

Platform: hotdocx.github.io

Sponsor: github.com/sponsors/hotdocx