

## Modelling Human Threats in Socio-Technical Systems

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Joint work with Giampaolo Bella and Carsten Schuermann

# Context

- Sociotechnical System

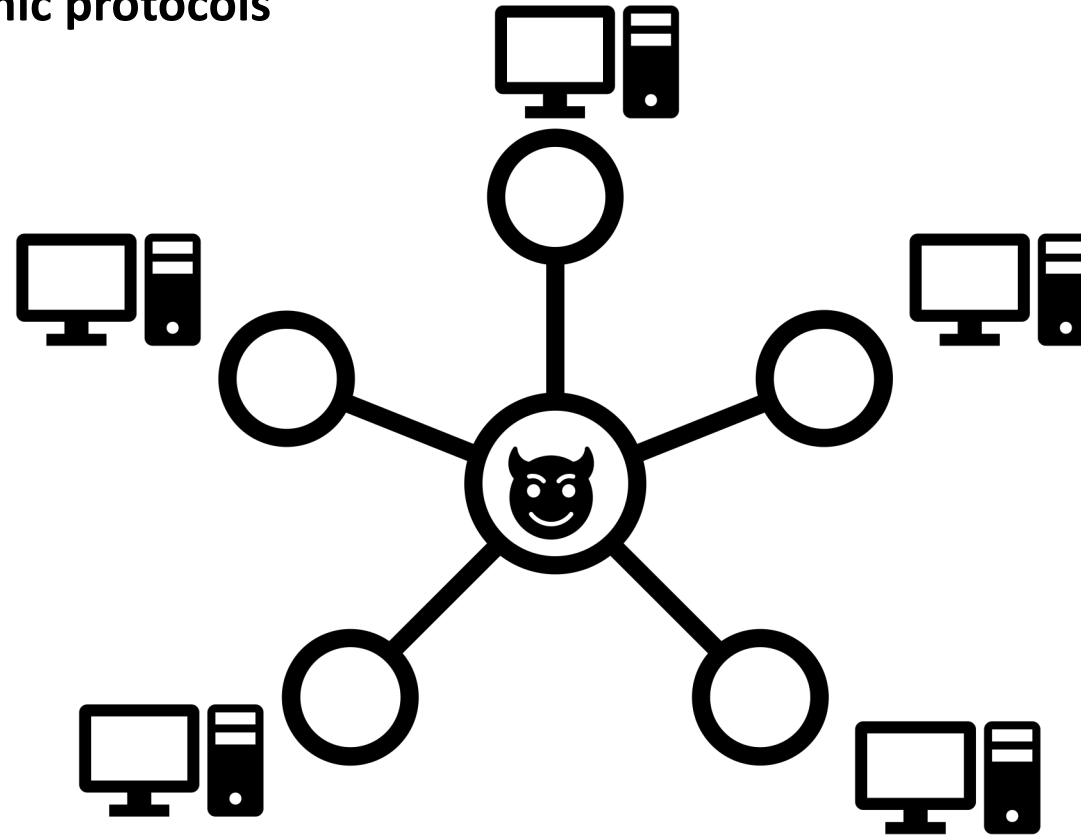
- A sociotechnical system is the term usually given to any instantiation of socio and technical elements engaged in **goal directed behaviour**. (*Wikipedia*)
- A technical system extended with its **human** users.
- **Security Ceremony**: A security **protocol** extended with its human users.

- Examples

- Flight boarding, **safety**
- Voting, **complex**
- POS transaction, **security**

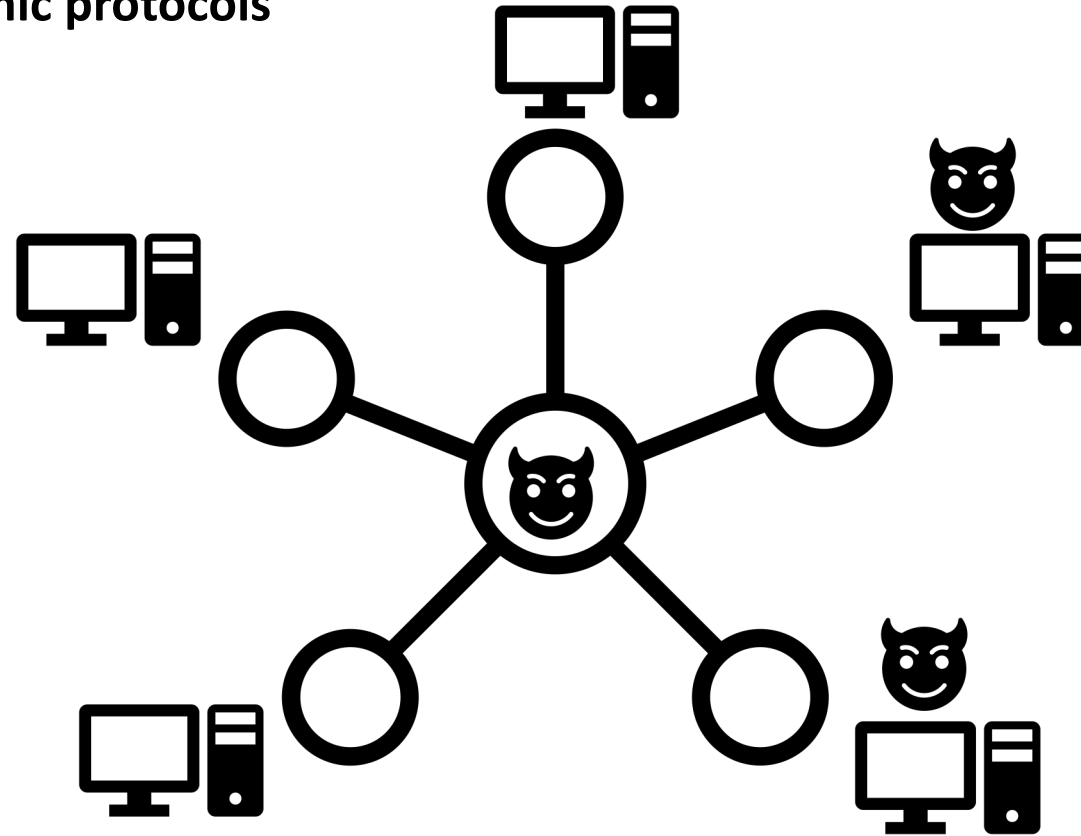
# “Not a problem”

- (formal) security analysis of
  - **Cryptographic protocols**



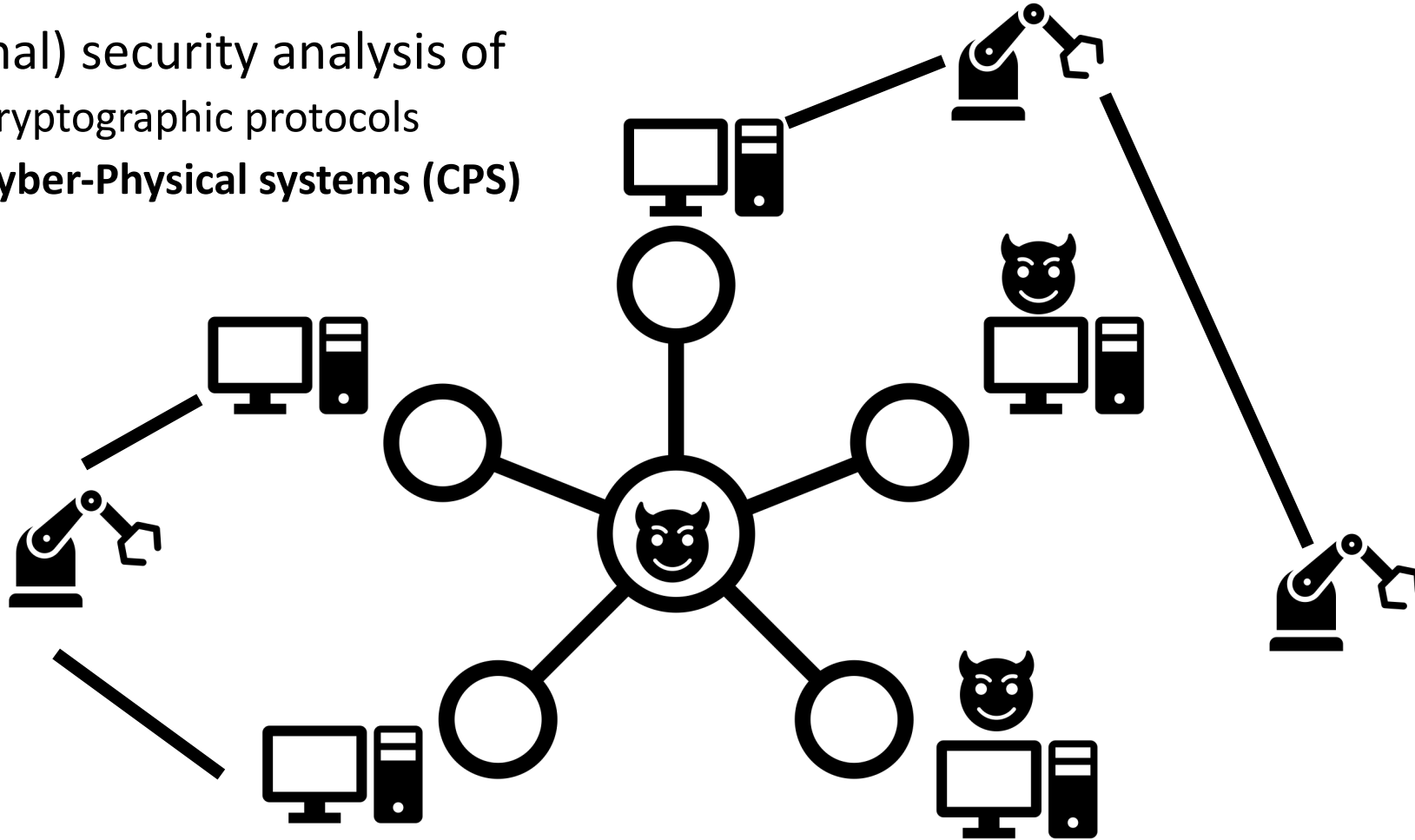
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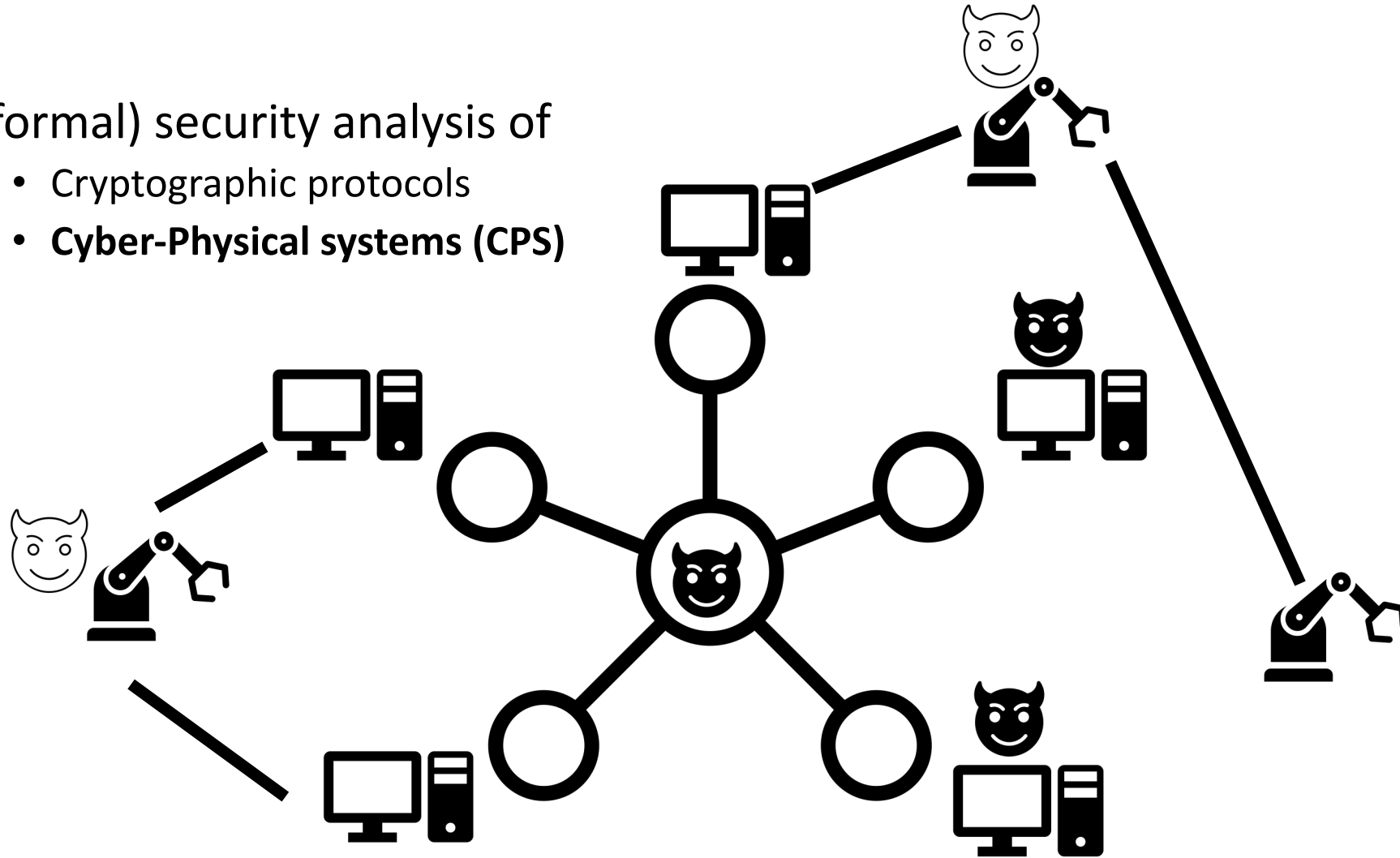
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- (formal) security analysis of
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  - **Cyber-Physical systems (CPS)**



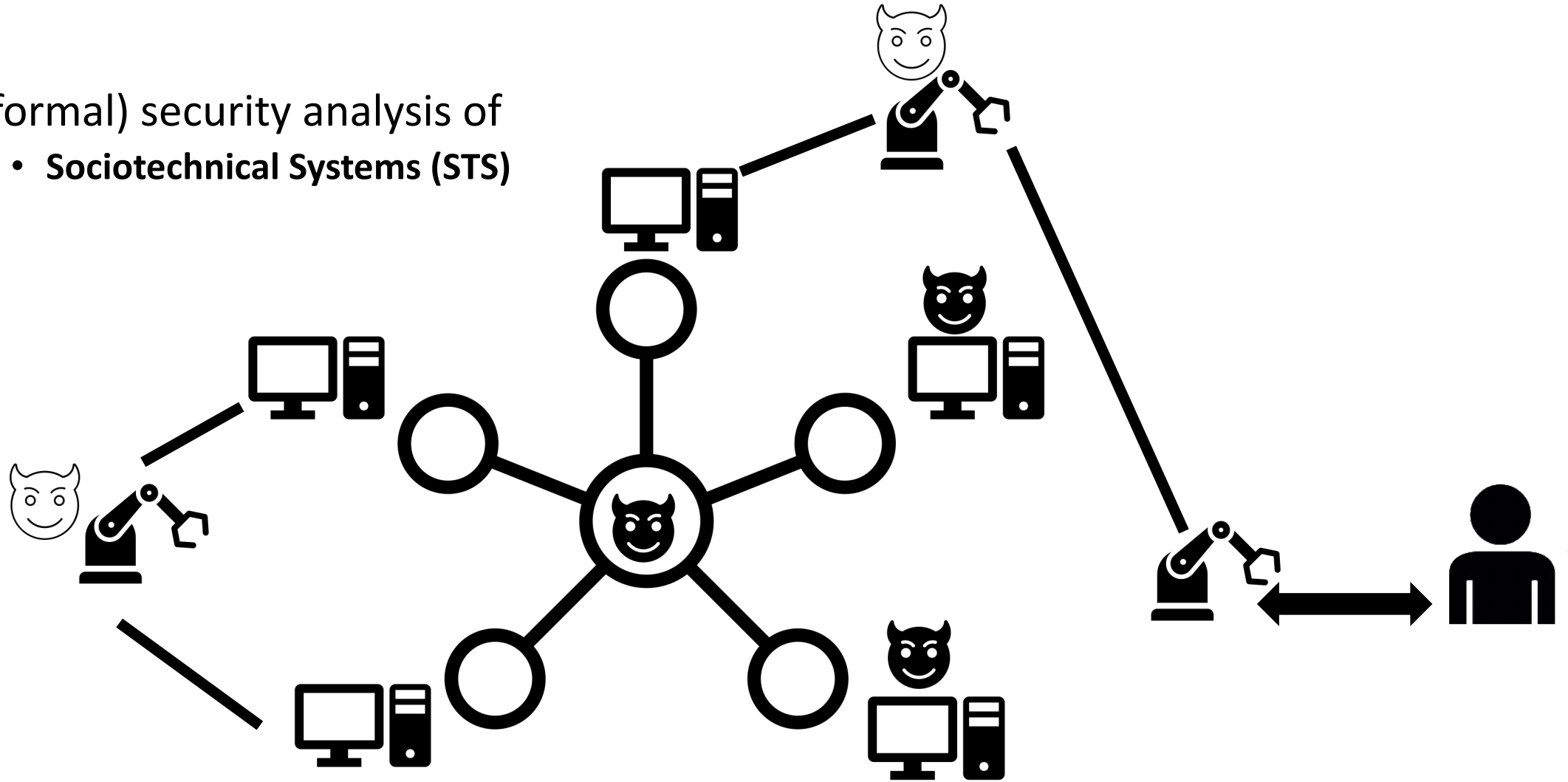
# “Not a problem”

- (formal) security analysis of
  - Cryptographic protocols
  - **Cyber-Physical systems (CPS)**



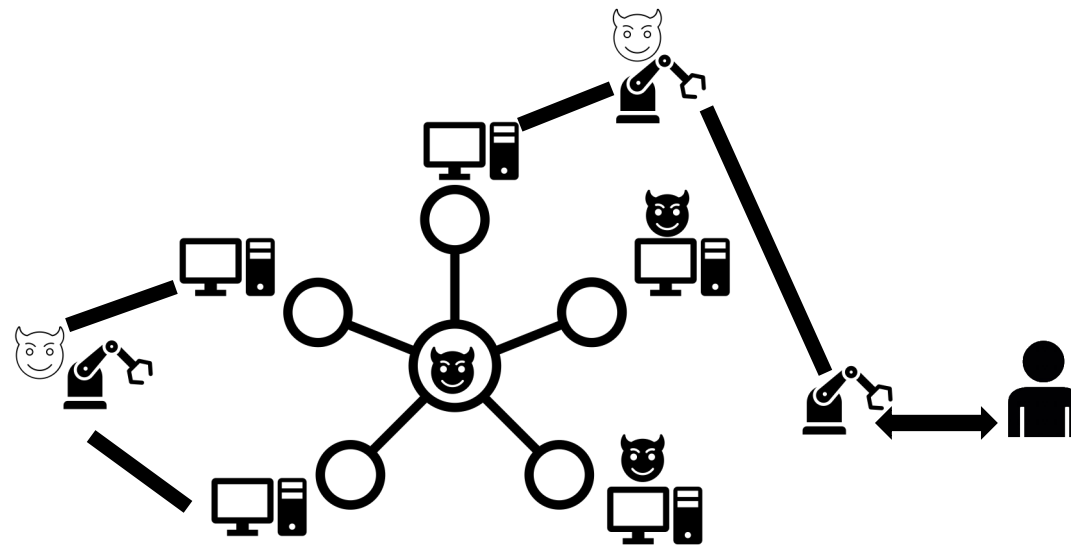
# Problem

- (formal) security analysis of
  - **Sociotechnical Systems (STS)**



# Problem

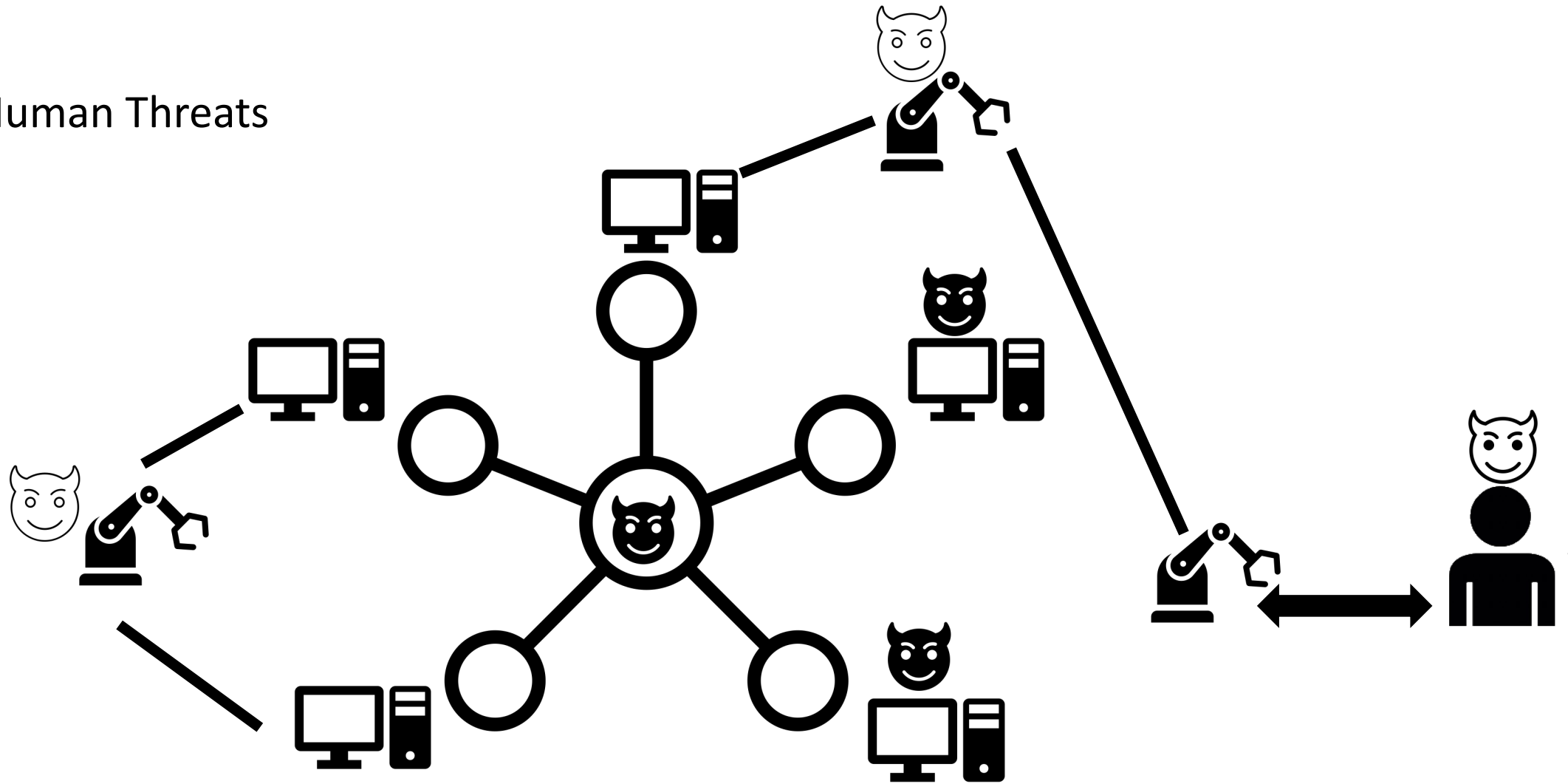
- Few works focus on the formal security analysis of STS
  - Bella and Coles-Kamp [IFIPSEC12], [focus on human-computer interaction](#)
  - Basin et al.[CSF16], [focus on human errors](#)
  - Sempereboni and Vigano [EuroSP20], [focus on “mutations” of humans and of the underlying technical components](#)





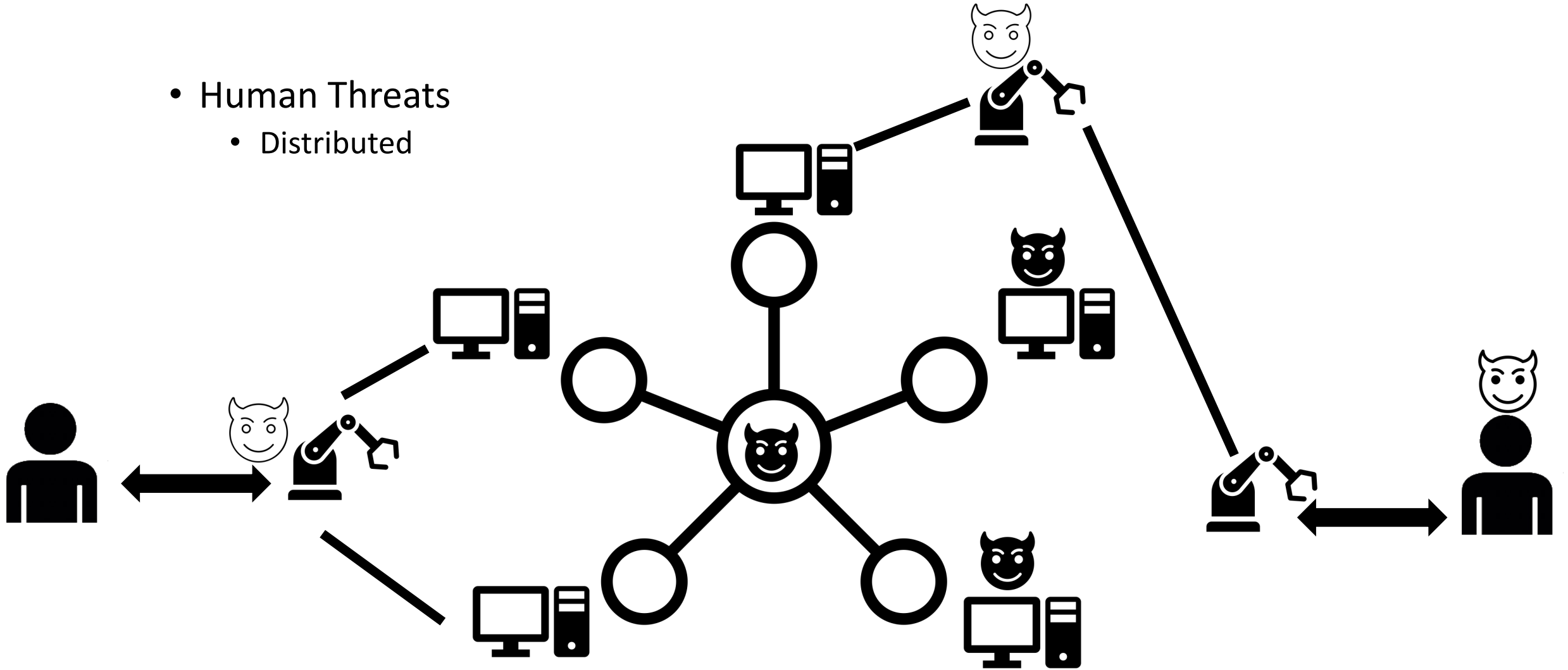
# Focus

- Human Threats



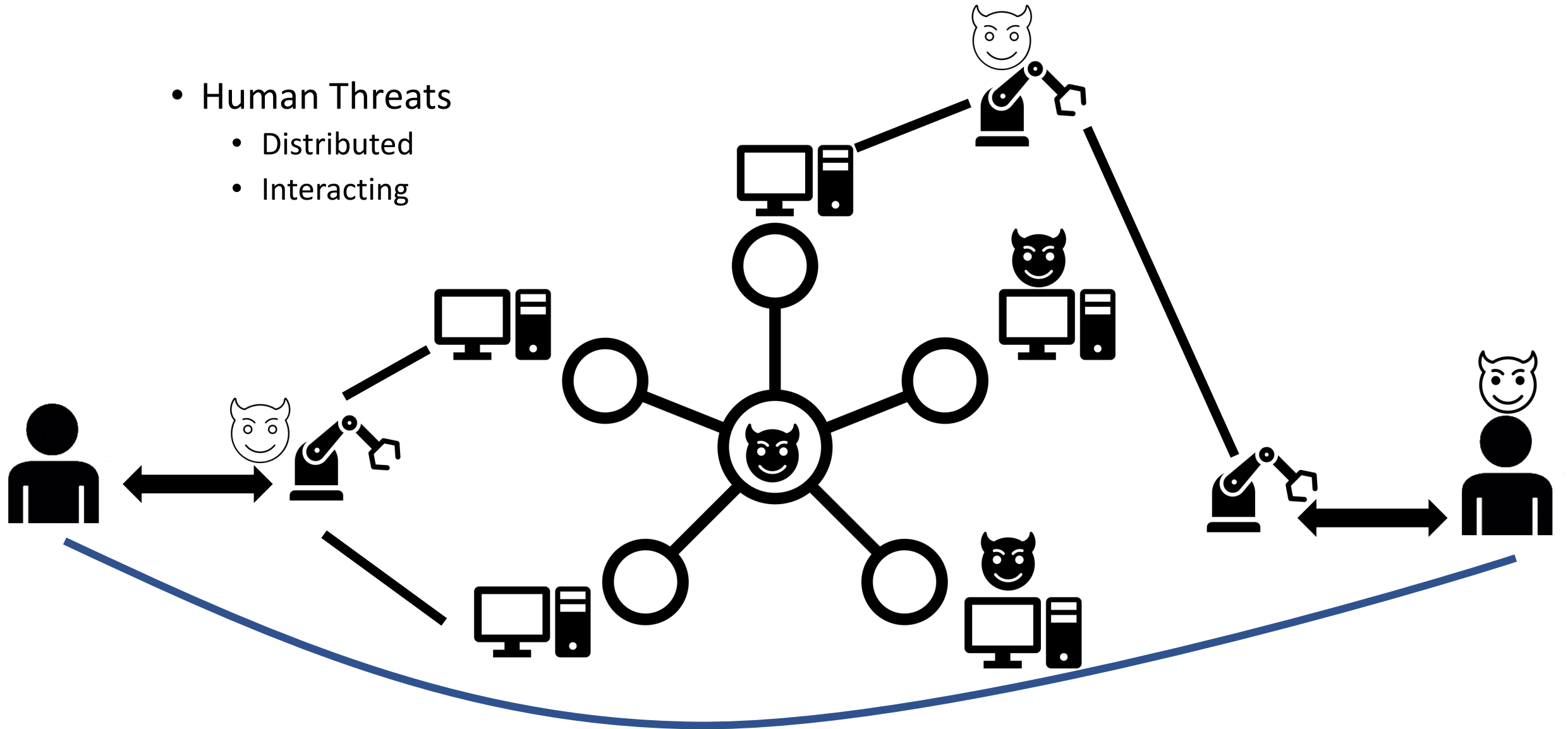
# Focus

- Human Threats
  - Distributed



# Focus

- Human Threats
  - Distributed
  - Interacting



# Outline

- Definitions in **epistemic modal logic of human threats** to STS
- Formal analysis of (Danish) **Deposit Return Systems** in Tamarin
- Definition of a **lattice of human threat** models
- Search methodology for finding **maximal threat models** not breaking security properties

# A threat model for interacting humans

- Essentially
  - *Honest*, **follows** the rules of a given ceremony precisely
  - *Chatty*, discloses their own **information**
    - e.g. reveals their passwords or the content printed on a ticket.
  - *Cocky*, gives out own **objects** that are relevant to the given ceremony
    - e.g., hands a physical token or a paper ticket.
  - *Forger*, **counterfeits** objects out of known information about them
    - e.g., builds physical token (provided they know the crypto material) via a 3D printer
    - *Receipt forger*, counterfeits printouts out of known information

# Epistemic modal logic

Terms  $t ::= x \mid f(t_1, \dots, t_n)$

Formulas  $F, G ::= 1 \mid P(t_1, \dots, t_n) \mid F \multimap G \mid F \otimes G \mid \llbracket K \rrbracket F \mid [K]F$   
 $\mid \forall x. F \mid \exists x. F \mid \Pi K. F \mid \Sigma K. F$

*(Look)* :  $\Pi P. \forall o. [P]object(o) \multimap (\llbracket P \rrbracket info(o) \otimes [P]object(o))$

# Epistemic modal logic

$(Chatty) : \Pi P. \Pi Q. \forall o. \llbracket P \rrbracket chatty \otimes \llbracket P \rrbracket info(o) \multimap \llbracket Q \rrbracket info(o)$

$(Give) : \Pi P. \Pi Q. \forall o. \llbracket P \rrbracket cocky \otimes [P]object(o) \multimap [Q]object(o)$

$(Hand) : \Pi P. \Pi Q, \forall r. \llbracket P \rrbracket cocky \otimes [P]receipt(r) \multimap [Q]receipt(r)$

$(Print) : \Pi K. \forall o. \llbracket K \rrbracket Rforger \otimes \llbracket K \rrbracket info(o) \multimap [K]receipt(QR(o))$

$(Build) : \Pi K. \forall b. \llbracket K \rrbracket Oforger \otimes \llbracket K \rrbracket info(b) \multimap [K]object(b)$

# Deposit Return System



Kæmpe boom i antallet af pantflasker:  
Hele 1,7 mia. flasker og dåser blev  
afleveret i 2020

22.3.2021 12:35:29 CET | [Dansk Retursystem](https://www.danskretur.dk)






# Problem

- Security protocols models are normally available
  - E.g. RFC, open-source, reverse engineering, etc.
- Sociotechnical system models are normally **not** available

## How to Board a Plane

 <http://www.wikihow.com/Board-a-Plane>

### Navigating the Airport

1. Print your boarding pass and check your luggage.
2. Head to security.
3. Find your gate/terminal.
4. Hang out and wait for your plane.



### Boarding the Plane

1. Wait for the announcement to board.
2. Get your boarding pass checked.
3. Walk down the hallway that leads up to your plane.
4. Enter the aircraft.
5. Stow your carry-on items.
6. Get settled in.

# Problem

- Security protocols models are normally available
  - E.g. RFC, open-source, reverse engineering, etc.
- Sociotechnical system models are normally **not** available



# Solution

- Field observation
  - To collect human behaviour
- Patents' analysis
  - To understand the technicalities
- Playing detective
  - To refine drafts of the ceremony

Customer



Cashier

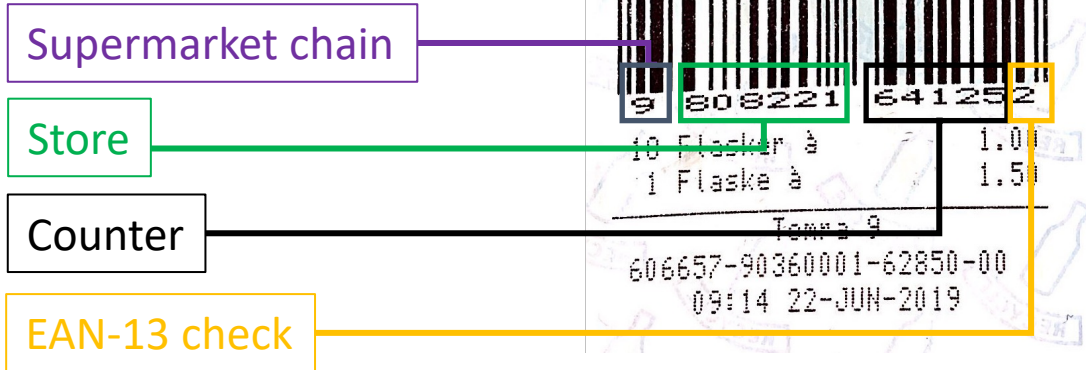


Return vending  
machine



Cash register

# Playing detective



# Playing detective

DENNE BON KAN  
KUN INDLØSES  
I KVICKLY  
NØRREBOGADE  
**Kvickly**

**4.00**  
2 339900 002004  
4 Flasker à 1.00  
Tomra 710  
606657-90159045-01273-00  
15:47 08-NOV-2018

Supermarket chain

Store

Counter

EAN-13 check

**Kvickly**  
**1.00**  
2 339900 001007  
1 Flaske à 1.00  
Tomra 710  
606657-90159045-01274-00  
16:47 07-NOV-2018

**Kvickly**  
**1.00**  
2 339900 001007  
1 Flaske à 1.00  
Tomra 710  
606657-90159045-01273-00  
16:07 06-NOV-2018

**Kvickly**  
**1.00**  
2 339900 001007  
1 Flaske à 1.00  
Tomra 710  
606657-90159001-01274-00  
14:12 07-FEB-2019

**Kvickly**  
**2.00**  
2 339900 002004  
2 Flasker à 1.00  
Tomra 710  
606657-90159045-01273-00  
15:47 08-NOV-2018

**Kvickly**  
**3.00**  
2 339900 003001  
3 Flasker à 1.00  
Tomra 710  
606657-90159045-01273-00  
19:56 09-NOV-2018

**Kvickly**  
**4.00**  
2 339900 004008  
4 Flasker à 1.00  
Tomra 710  
606657-90159045-01273-00  
11:44 17-NOV-2018

**coop**  
**2.00**  
2 139900 002000  
2 Flasker à 1.00  
Tomra 710  
606657-90159045-06949-00  
14:38 07-FEB-2019

**coop**  
**3.00**  
2 139900 003007  
1 Flaske à 0.00  
3 Flasker à 1.00  
Tomra 710  
606657-90159045-06949-00  
14:39 07-FEB-2019

**coop**  
**3.00**  
2 139900 003007  
2 Flasker à 1.50  
Tomra 710  
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14:37 07-FEB-2019

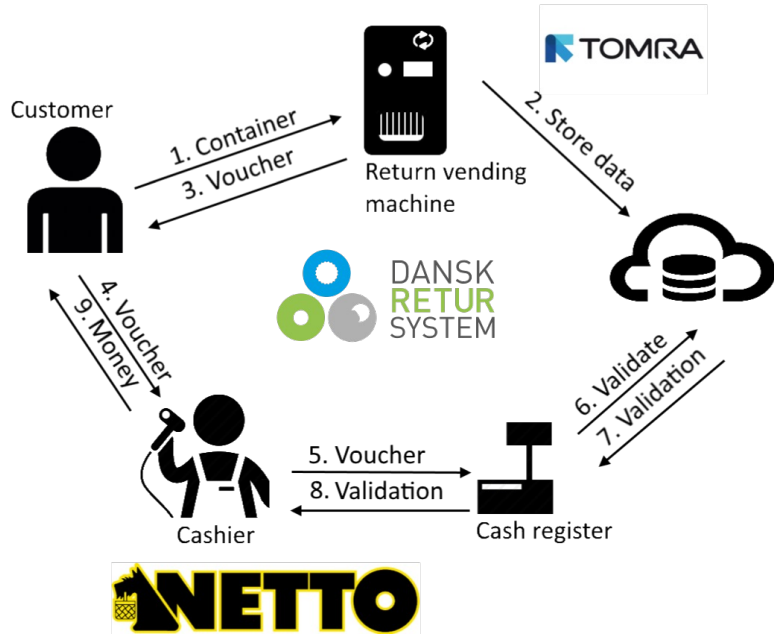
KONGELUNDSVEJ 49.  
**NETTO**  
**kr 19.50**  
9 808221 641252  
18 Flasker à 1.00  
1 Flaske à 1.50  
Tomra 9  
606657-90360001-62850-00  
09:14 22-JUN-2019

KONGELUNDSVEJ 49.  
**NETTO**  
**kr 20.50**  
9 808221 641245  
1 Flaske à 1.00  
3 Flasker à 1.50  
5 Flasker à 3.00  
Tomra 9  
606657-90360001-62850-00  
09:11 22-JUN-2019

KONGELUNDSVEJ 49.  
**NETTO**  
**kr 11.00**  
9 808221 641269  
11 Flasker à 1.00  
Tomra 9  
606657-90360001-62850-00  
09:15 22-JUN-2019



# Deposit Return System



$(Purchase) : \Pi S. \Pi C. \forall c. \llbracket S \rrbracket seller \otimes \llbracket C \rrbracket customer \otimes [S] object(c) \multimap [C] object(c)$

$(Return) : \Pi C. \Pi V. \forall c. \llbracket C \rrbracket customer \otimes \llbracket V \rrbracket rvm \otimes [C] object(c) \multimap [V] object(c)$

$(Output) : \Pi V. \Pi Ca. \Pi C. \forall c. \llbracket V \rrbracket rvm \otimes \llbracket Ca \rrbracket cashier \otimes \llbracket C \rrbracket customer \otimes [V] object(c) \multimap \llbracket Ca \rrbracket info(c) \otimes \exists id. [C] receipt(QR(c, id))$

$(Hand) : \Pi C. \Pi Ca. \forall r. \llbracket C \rrbracket customer \otimes \llbracket Ca \rrbracket cashier \otimes [C] receipt(r) \multimap [Ca] receipt(r)$

$(Cash) : \Pi Ca. \forall id, \forall c. \llbracket Ca \rrbracket cashier \otimes [Ca] receipt(QR(c, id)) \otimes \llbracket Ca \rrbracket info(c) \multimap 1$

# Formal analysis

- Tamarin
  - Essentially a *constraint solver*
  - **Parties and threat specs using *multi-set rewriting***
  - **Properties spec using *metric first-order logic***
  - Proofs constructed using *backward search*

Encoding epistemic modal logic → Tamarin

# Encoding epistemic modal logic → Tamarin

- Properties as metric first-order logic

- E.g. forall traces, *Cash* is always preceded by a container *Hand*.

*(Purchase)* :  $\Pi S. \Pi C. \forall c. \llbracket S \rrbracket seller \otimes \llbracket C \rrbracket customer \otimes [S] object(c) \multimap [C] object(c)$

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$\forall Ca C id c \#i. Cash(Ca, id, c)@i \implies \exists \#j. Hand(C, Ca, QR(c, id))@j \wedge j < i$

# Properties

If a **cashier** cashes out a **voucher**, then a corresponding...

**Cash for voucher** ... **receipt** has been printed earlier by a **RVM**

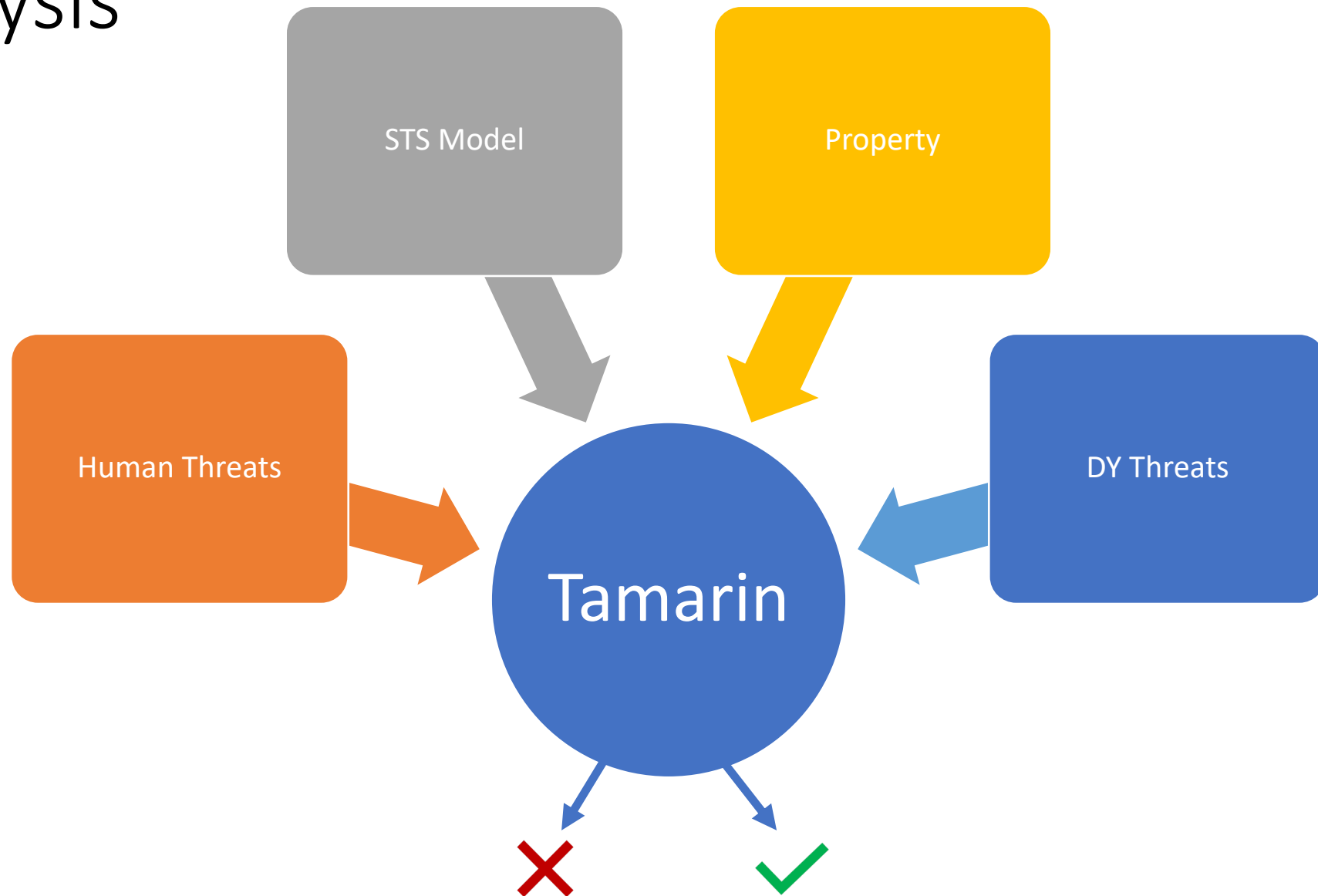
**Cash for container** ... **container** has been returned earlier to a **RVM**

**Strong cash for container** ... **container** has been returned earlier to a **RVM** by the **buyer**

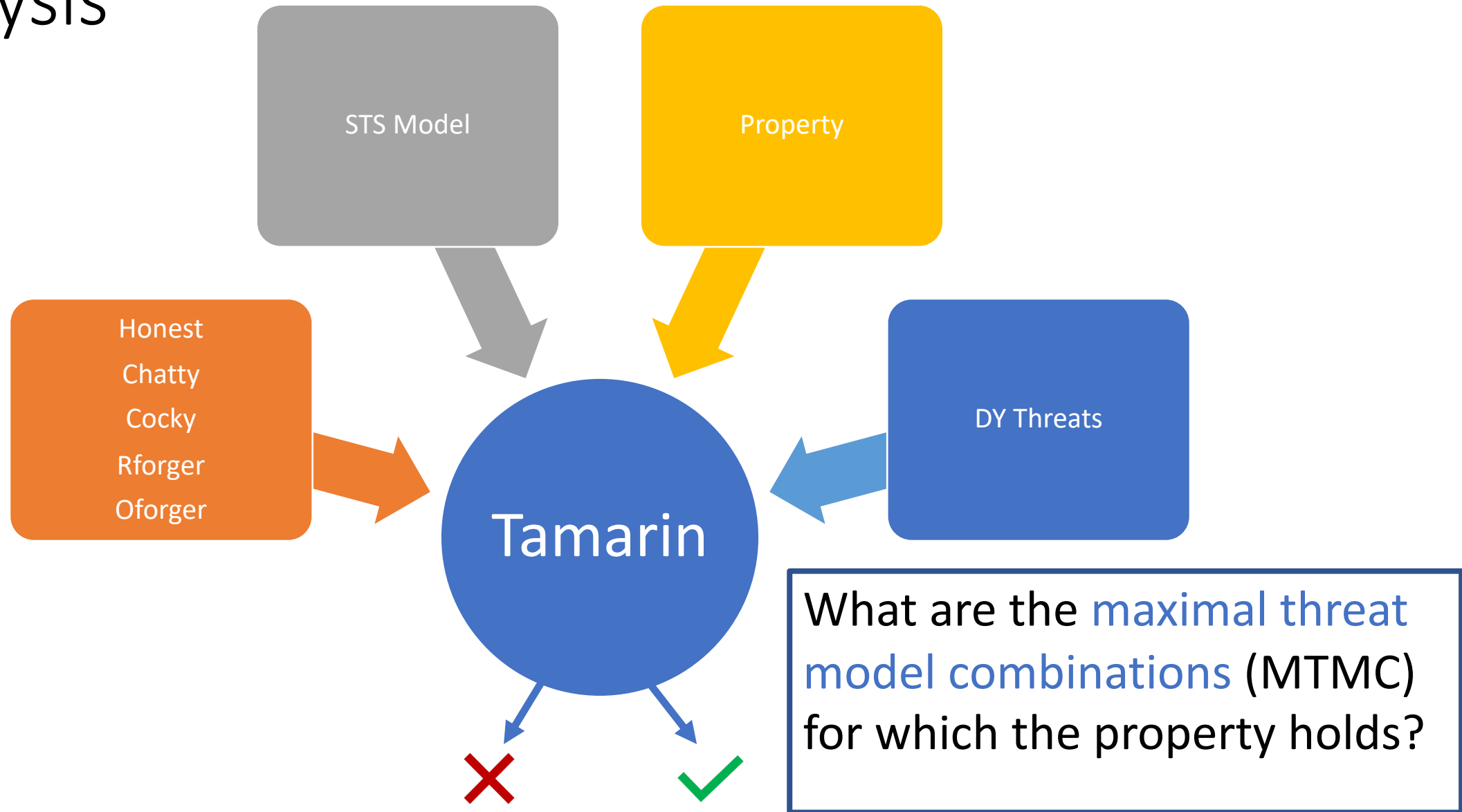
**Cash for purchase** ... **container** has been bought earlier

**Strong cash for purchase** ... **container** has been bought earlier by the same **customer**

# Analysis



# Analysis



# Lattice of human threats



$$Ch \wedge Co \wedge RF \wedge OF$$

# Lattice of human threats

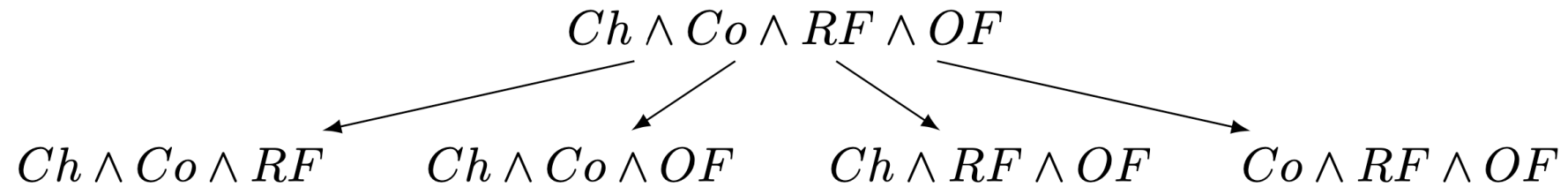
Honest

Chatty

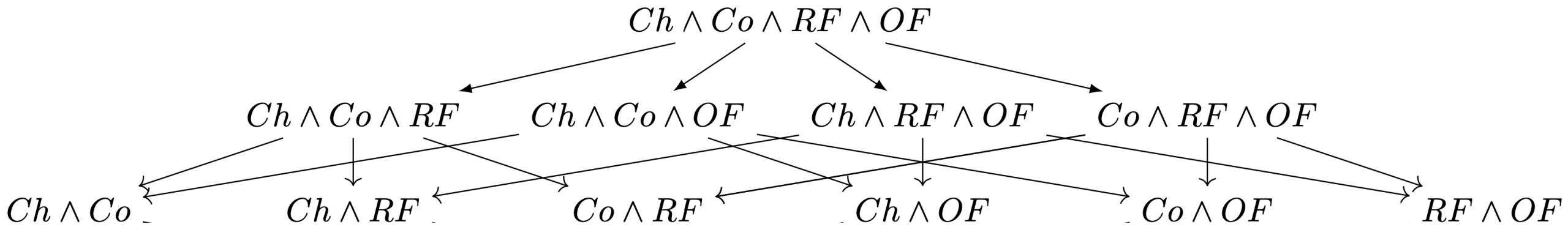
Cocky

Rforger

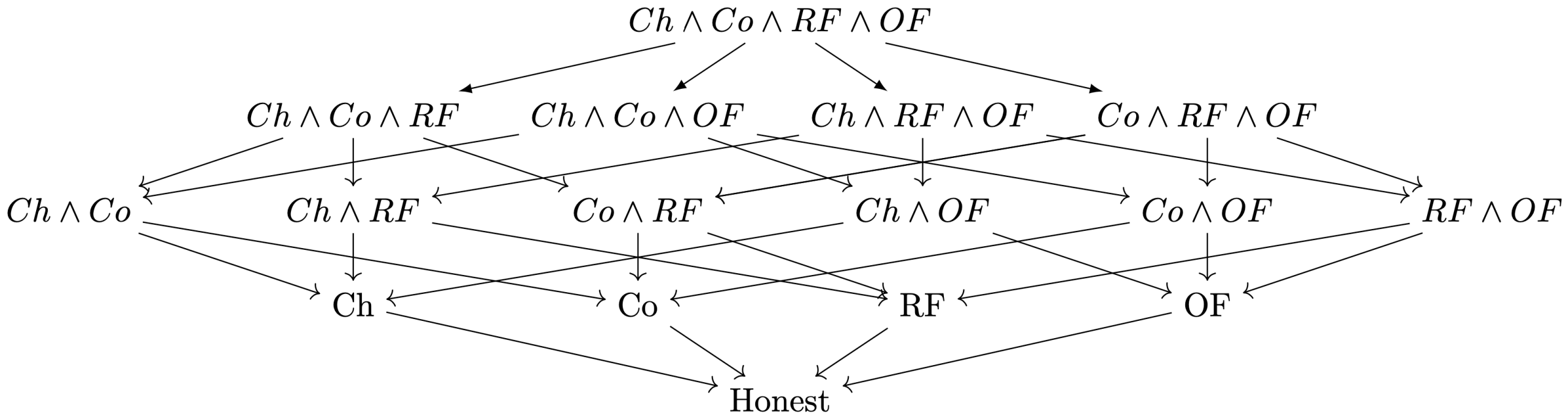
Oforger



# Lattice of human threats



# Lattice of human threats

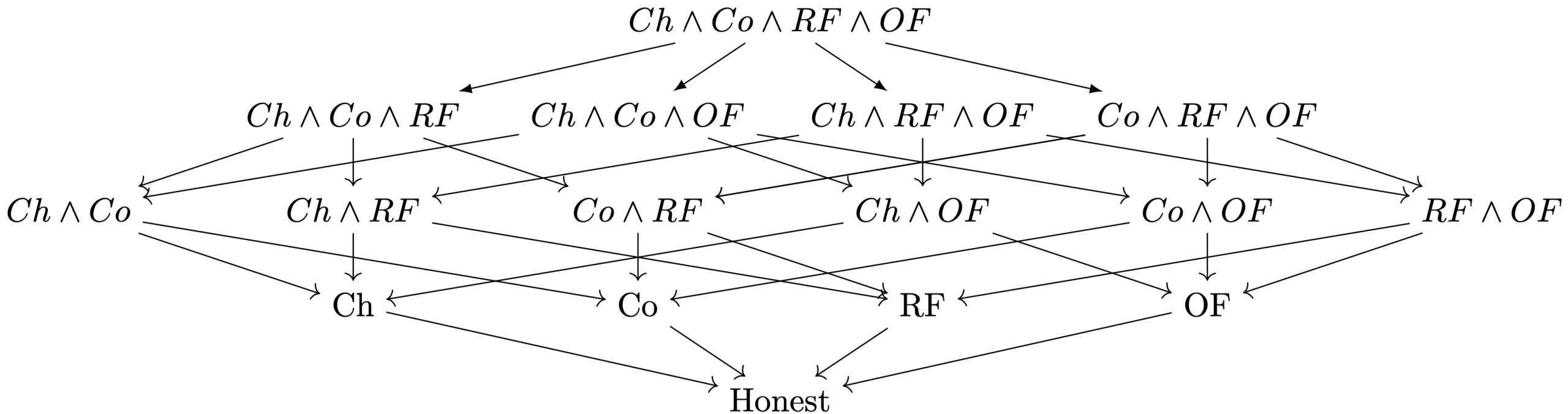




# Lattice of human threats

What are the maximal threat model combinations (MTMC) for which the property holds?

- Honest
- Chatty
- Cocky
- Rforger
- Oforger

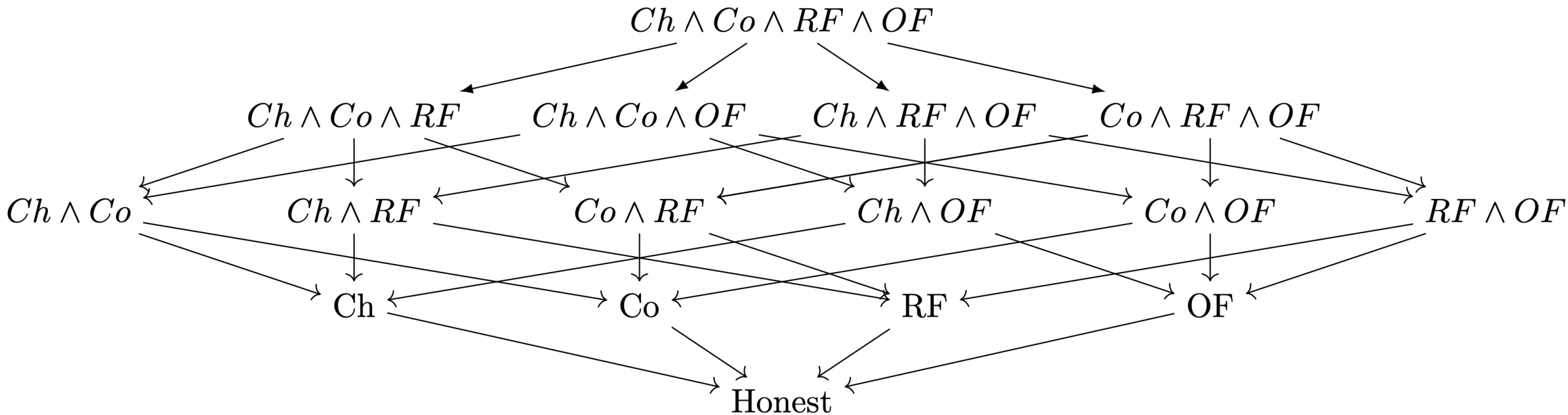


# Lattice of human threats

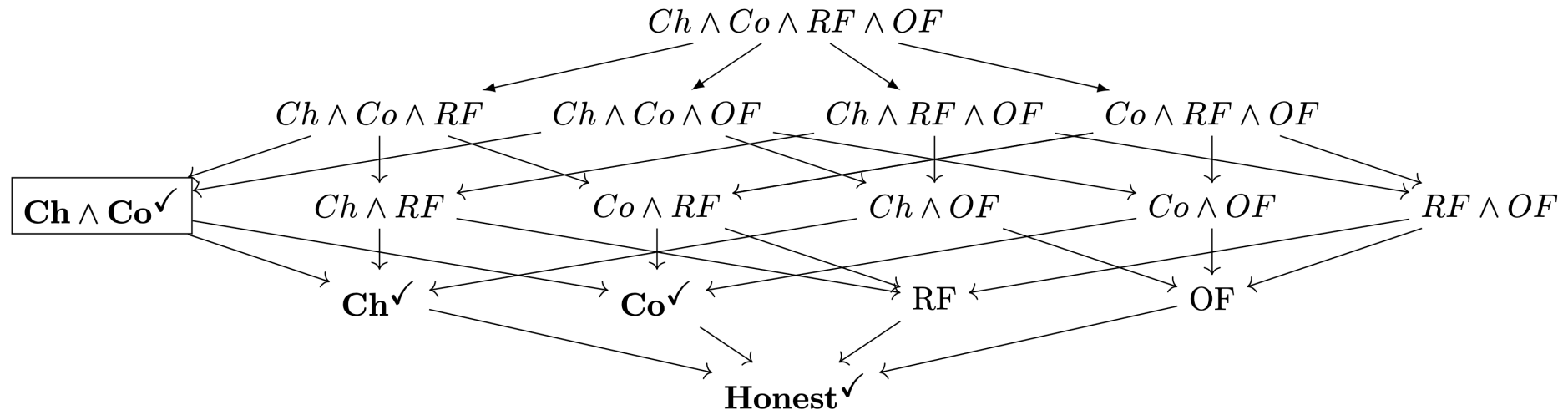
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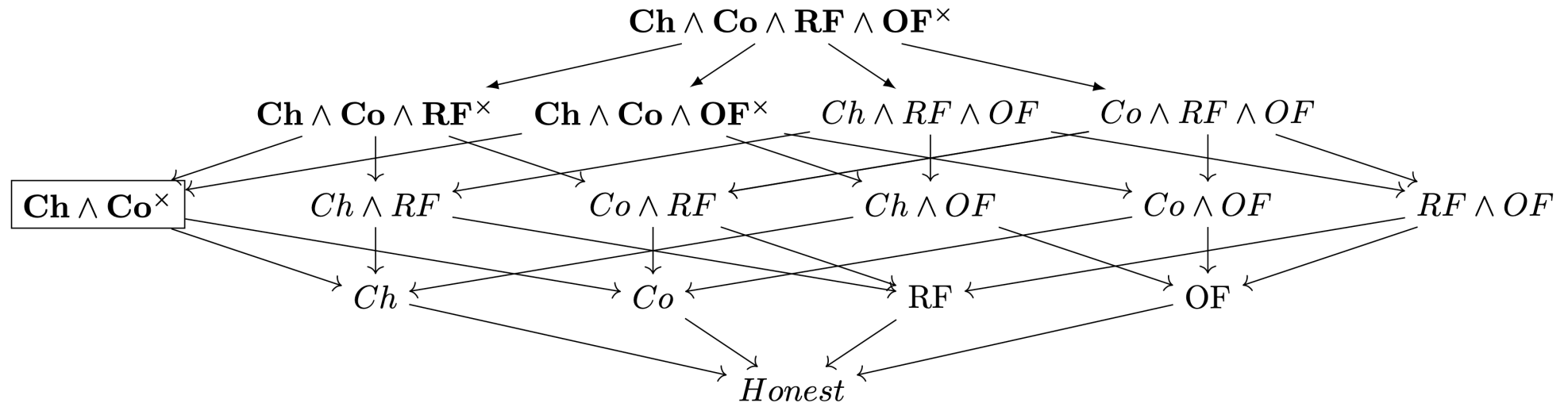
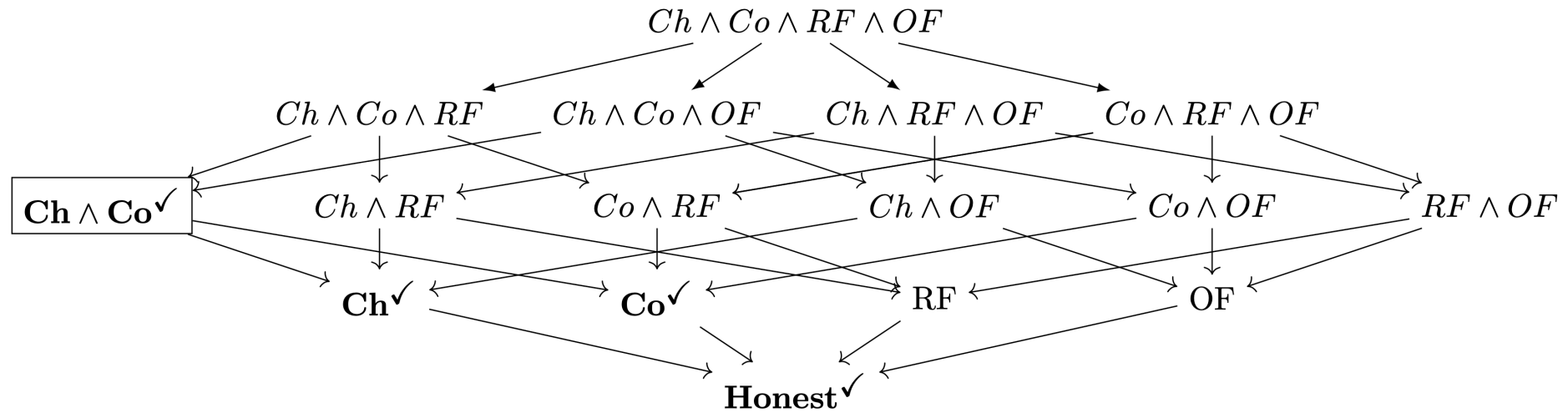
Problem: Finding MTMC can be time expensive!  
Worst case:  $2^{\text{#threats}}$  Tamarin analysis per property



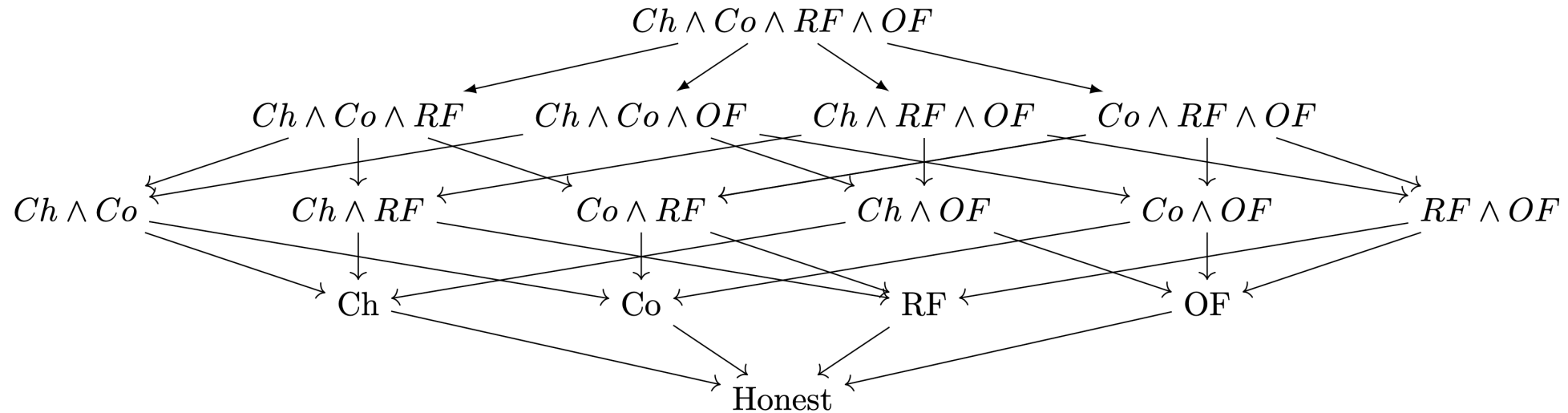
# Lattice of human threats



# Lattice of human threats



# Lattice of human threats



**procedure** Greedy\_check ( $G$ ):

$c \leftarrow \text{Max\_edges}(G)$

**if**  $L(c) \rightarrow \checkmark$  **then**

  | DFS ( $G, c, \checkmark, \downarrow$ )

**else**

  | DFS ( $G, c, \times, \uparrow$ )

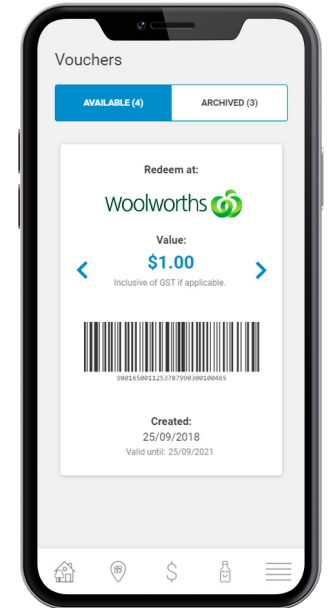
Greedy\_check ( $G [C - C_l]$ )

# Results

	Kvickly & Coop		Netto	
	<i>Result</i>	<i>MTMC</i>	<i>Result</i>	<i>MTMC</i>
Cash for voucher	×	$(Ch \wedge Co \wedge OF), (RF \wedge OF)$	✓	$(Ch \wedge Co \wedge RF \wedge OF)$
Cash for container	×	$(Ch \wedge Co \wedge OF), (RF \wedge OF)$	✓	$(Ch \wedge Co \wedge RF \wedge OF)$
Cash for container customer	×	$(Ch \wedge OF), (RF \wedge OF)$	×	$(Ch \wedge OF), (RF \wedge OF)$
Cash for purchase	×	$(Ch \wedge Co), (RF \wedge OF)$	×	$(Ch \wedge RF \wedge Co), (RF \wedge OF)$
Cash for purchase customer	×	$Ch, (RF \wedge OF)$	×	$Ch, (RF \wedge OF)$

# Fix

- Inspiration from *myTomraApp*, piloted in Australia
  - Cash out directly at the RVM







# Conclusion

- Attempt to **understanding formally** human threats in STS
- **Lattice of threat models** makes sense when dealing with human threats
- Fixes against human threats require a **shifting to technical solutions**
- Just an attempt
  - A general set of human + physical + network threats
  - Encoding epistemic modal logic to Tamarin
  - More case studies
  - Consider privacy

*“under the Danish law it is not allowed to copy or make changes to vouchers or to encourage others to do so.”*